`dynamic_cast` from scratch

Two kinds of inheritance graphs

Arthur O’Dwyer
2017-09-26

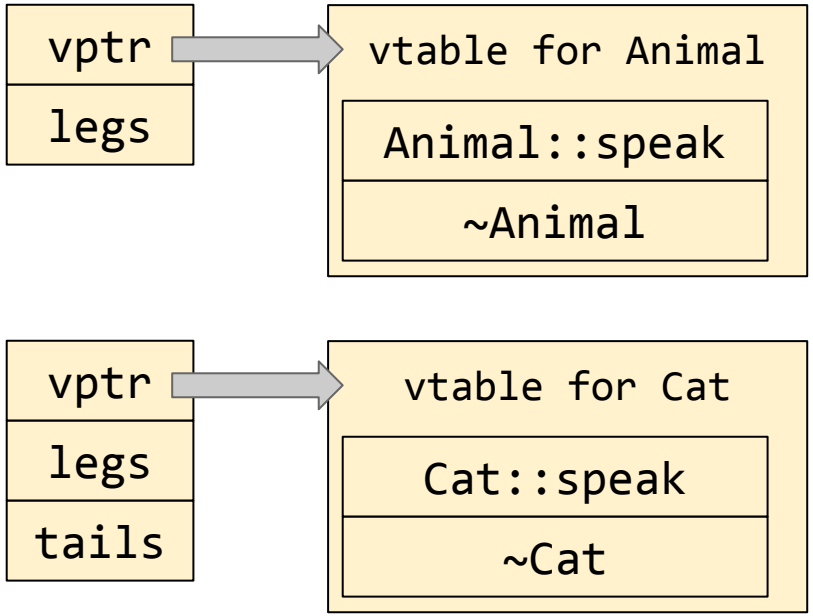
Outline

- **Part I: Recap polymorphism**
  - Inheritance, vtables, and the Dreaded Diamond [3–10]
  - Accessing members of virtual bases [11–21]
  - How to visualize complex class layouts [22–29]
- Part II: `dynamic_cast`
  - What should `dynamic_cast` do? [31–54]
  - Okay, how do we implement that? [55–60]
  - Benchmark numbers [61–64]

Quick recap: Polymorphism

```
class Animal {
    int legs;
    virtual void speak() { puts("hi"); }
    virtual ~Animal();
};

class Cat : public Animal {
    int tails;
    void speak() override {
        printf("Ouch, my %d tails!",
            tails);
    }
};
```



Quick recap: Polymorphism

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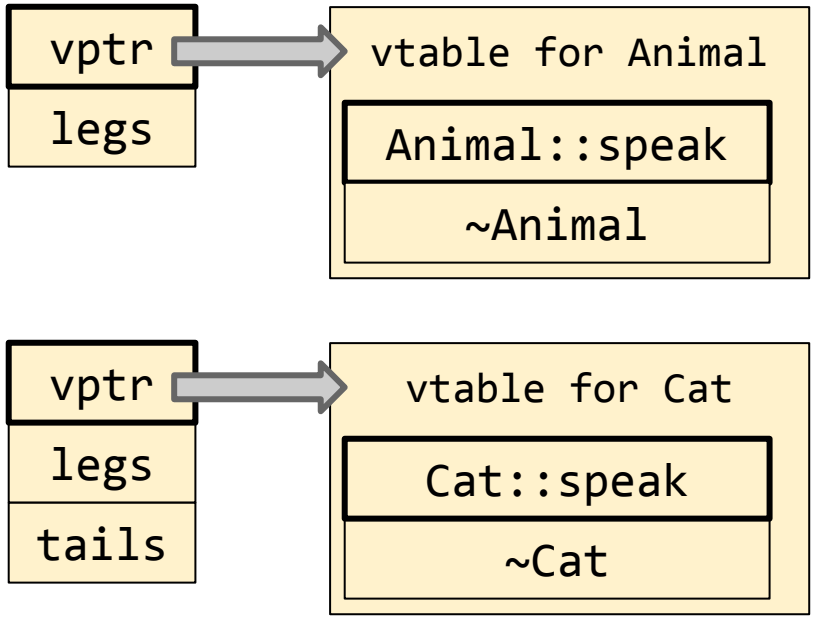
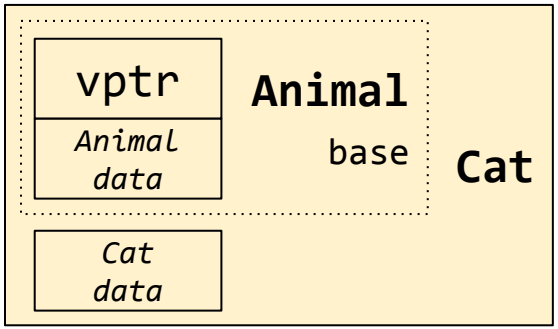
    void speak() override {
        printf("Ouch, my %d tails!",
            tails);
    }
};
```

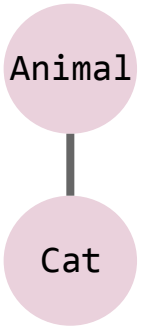
```
# a->speak();
movq  (%rdi), %rax
callq *(%rax)
```



Several graphical representations

Cat **IS-AN** Animal





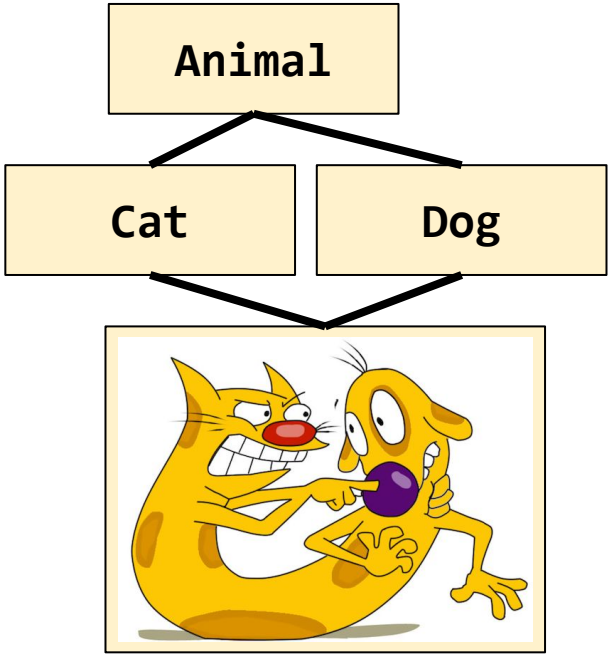
Multiple inheritance

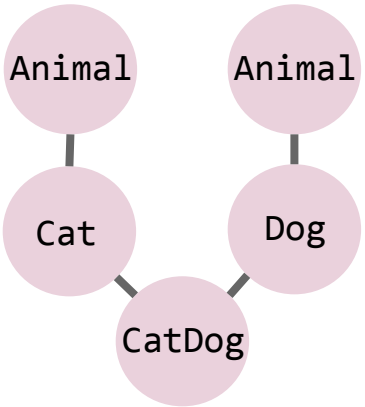
```
class Animal { virtual ~Animal(); };

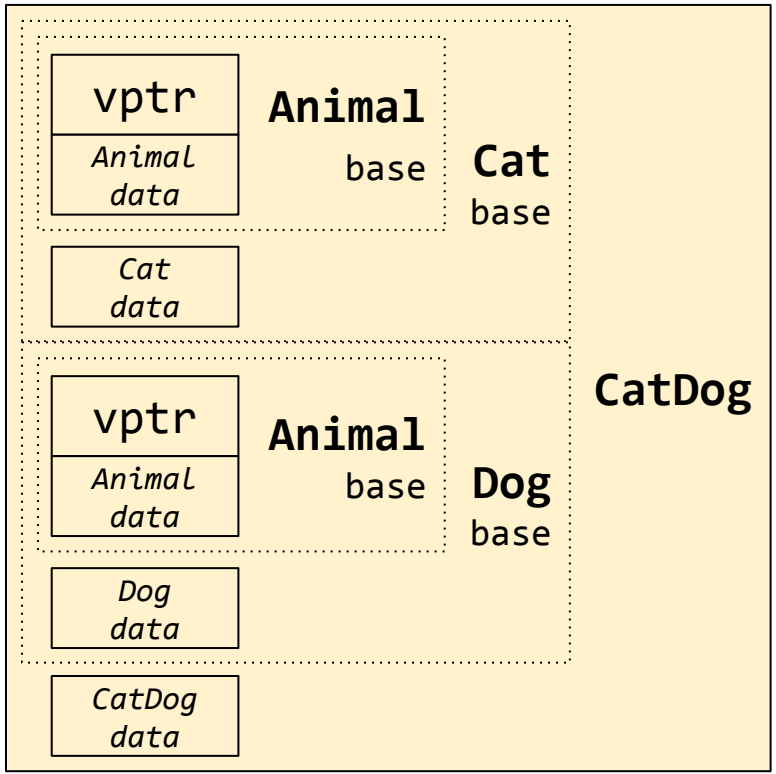
class Cat : public Animal { };

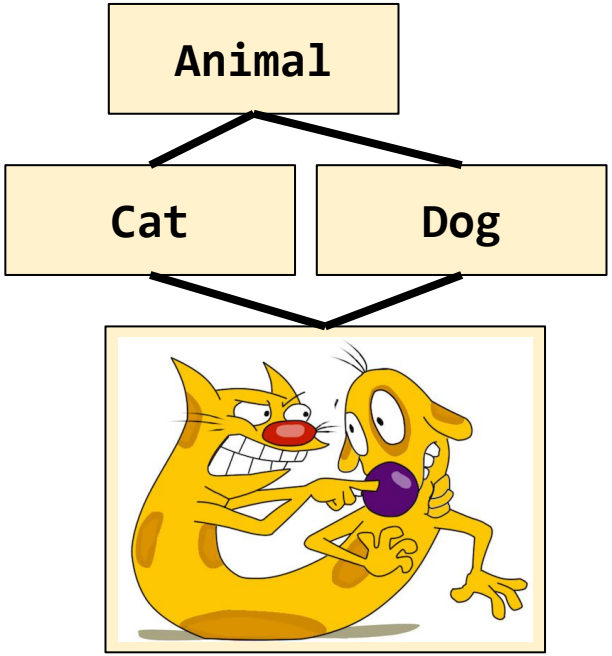
class Dog : public Animal { };


class CatDog :
  public Cat, public Dog { };
```

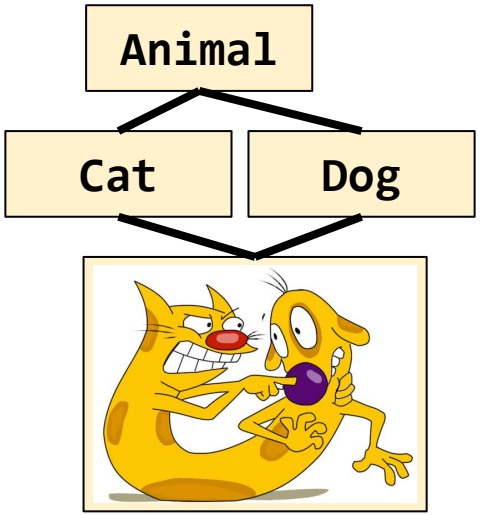



Animal
Animal
Cat
Dog
CatDog

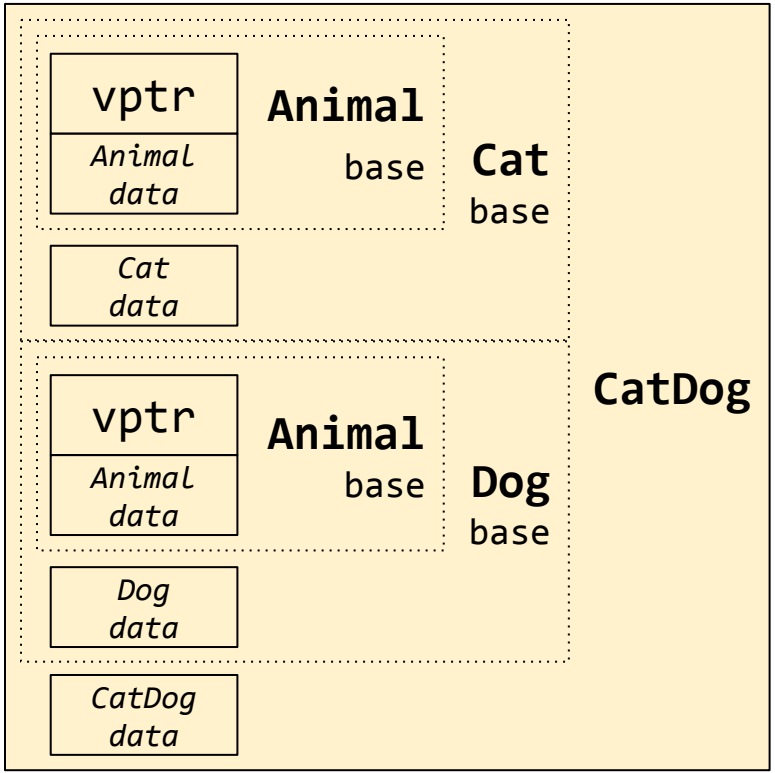

vptr
Animal data
Animal base
Cat base
Cat data
vptr
Animal data
Animal base
Dog base
Dog data
CatDog data
CatDog

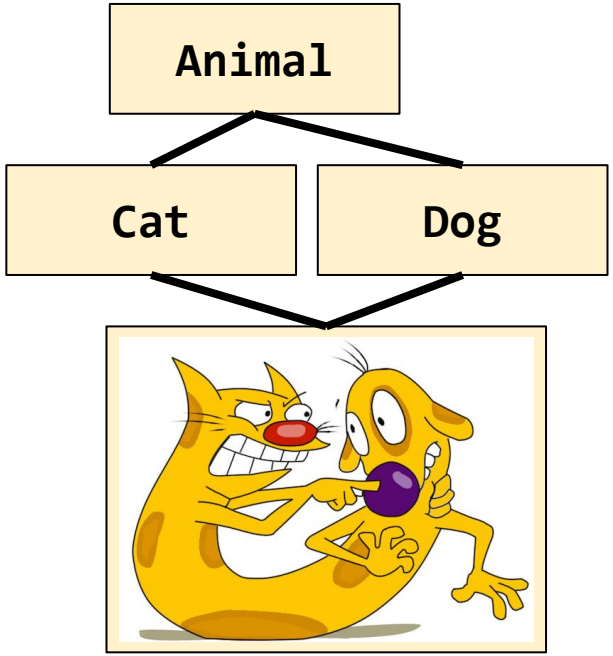

Animal
Cat
Dog

IS-A CatDog an Animal?

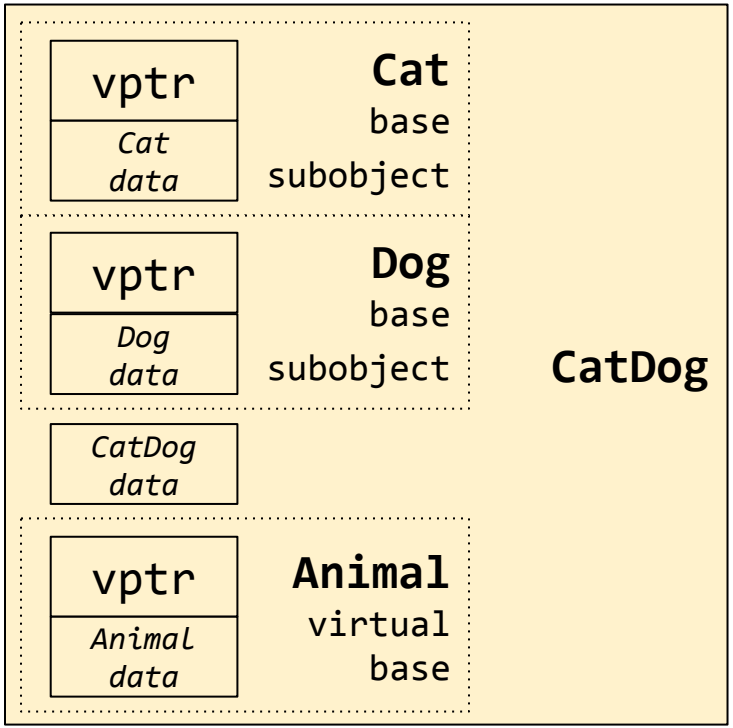


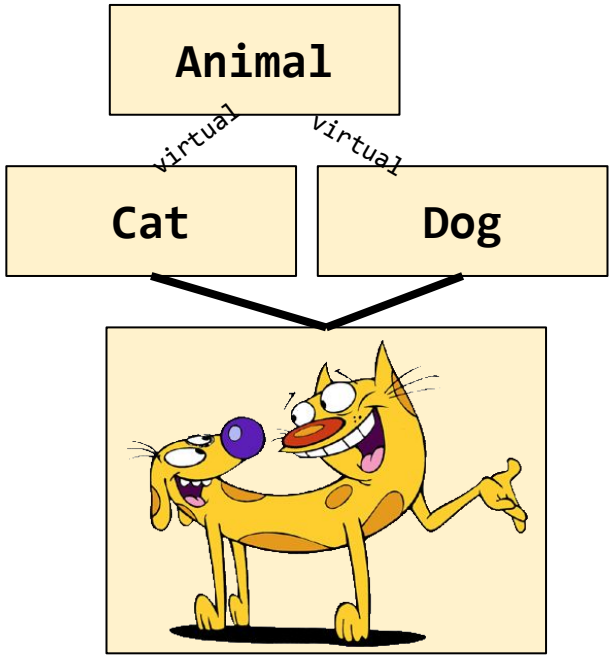
IS-A CatDog an Animal? No. (It's two Animals.)





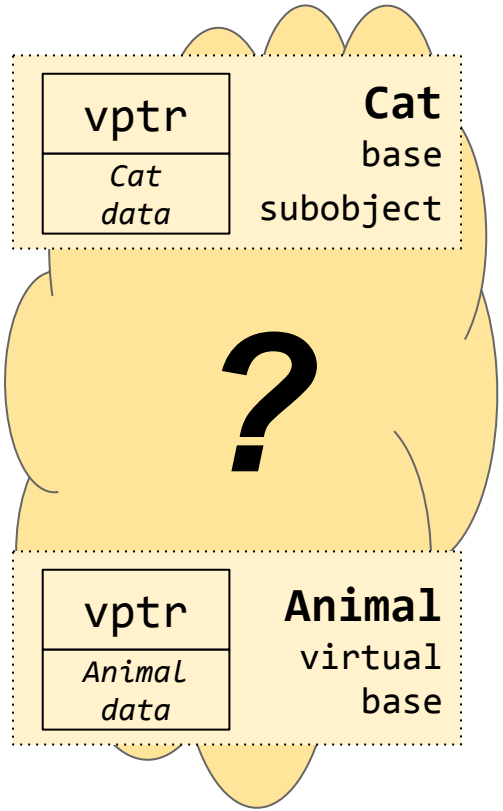
Fix the Dreaded Diamond with *virtual inheritance*

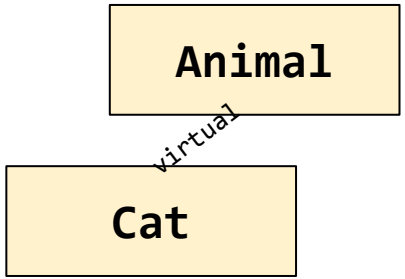




Now IS-A CatDog an Animal? Definitely yes.

But now we've changed how a Cat looks!

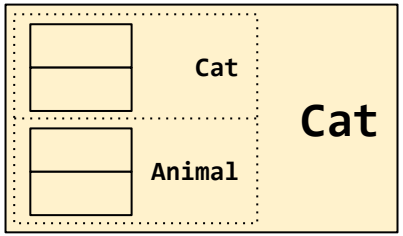




How do we get at the Animal data?

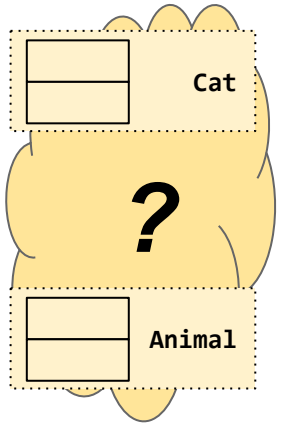
How do we get at the Animal data?

We have to distinguish between objects that are really just `Cats` and objects that are "`Cat` but maybe more."



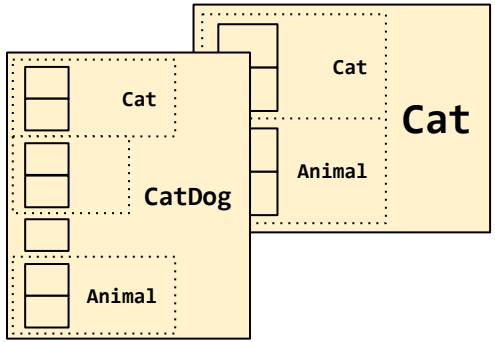
`Cat` objects that are **just `Cat`** have a fixed, concrete layout.

`Cat` objects that are **maybe more** have a fixed layout for all their "base subobject" pieces, but you have to do some work to figure out where the virtual bases are located.



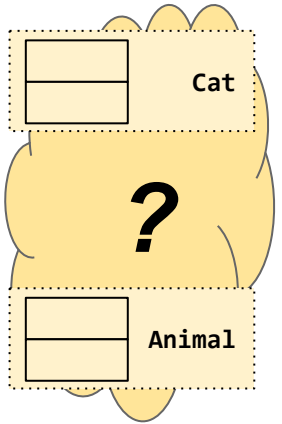
How do we get at the Animal data?

We have to distinguish between objects that are **most derived** and objects whose dynamic type is unknown.



`Cat` objects that are **most derived** have a fixed, concrete layout, dictated by their most derived type.

`Cat` objects of **unknown dynamic type** have a fixed layout for all their "base subobject" pieces, but to access a virtual base, you first have to figure out ***what is the most derived type of this object.*** That tells you its fixed, concrete layout; and *that* tells you where its virtual bases are located.



Vtables are always controlled by the most derived object

The *schema* of a vtable is dictated by the object's static type.

- A vtable "for `Cat` in (foo)" will always have a `speak()` method pointer at offset 0, a destructor at offset 8, and so on.
- A vtable "for `Dog` in (foo)" will always have a destructor at offset 0, and so on.

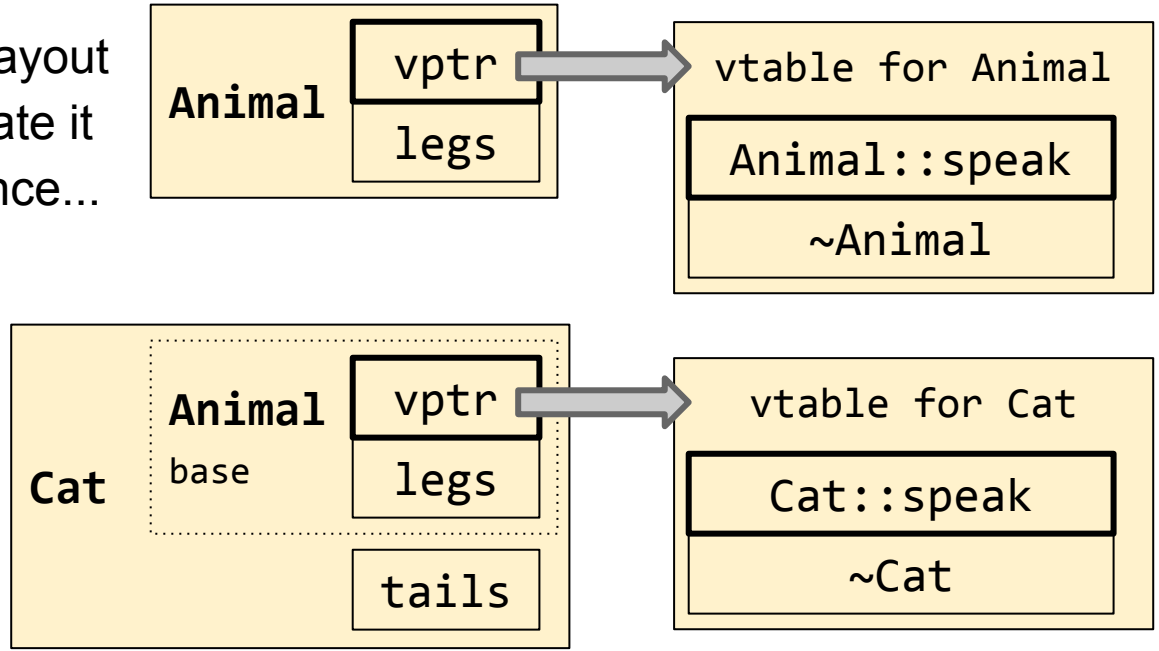
The *data* in a vtable is dictated by the object's dynamic (most-derived) type.

The *presence* of the virtual base `Animal` is dictated statically by the fact that the object **IS-A** `Cat`. The *location* of the virtual base `Animal` is dictated dynamically by the dynamic (most-derived) type.

We can find our virtual base `Animal` by querying our `Cat` object's vtable.

Vtables controlled by most-derived

This was how our memory layout looked on slide 4. Let's update it to reflect the virtual inheritance...



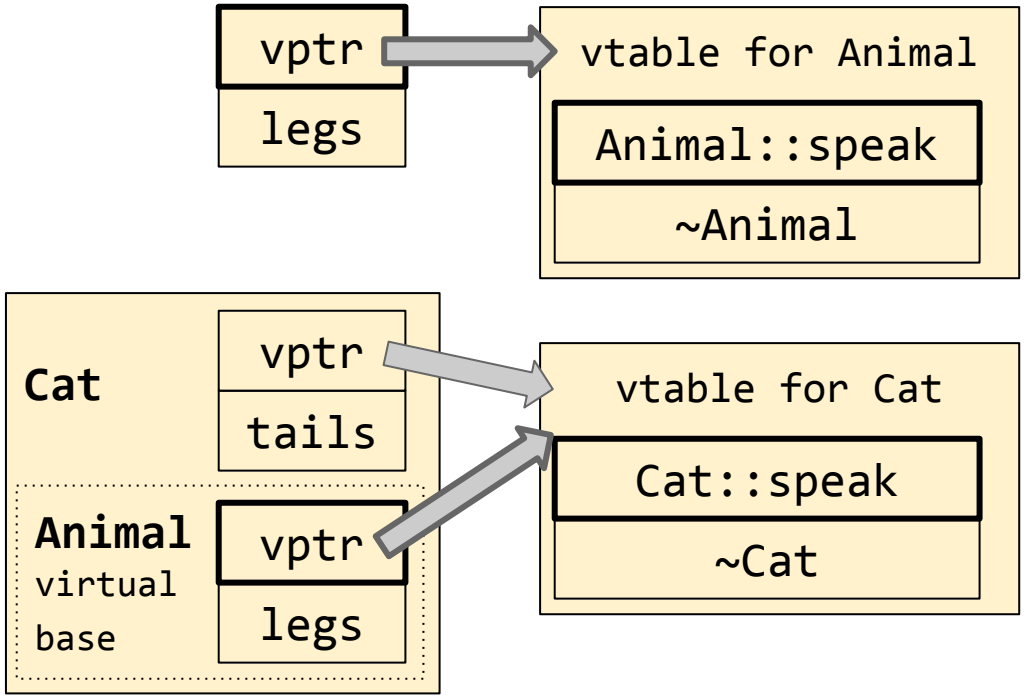
Notice that this is also the vtable for "Animal in Cat," so its schema is a superset of the vtable for "Animal in Animal" above.

Vtables controlled by most-derived

This was how our memory layout looked on slide 4. Let's update it to reflect the virtual inheritance...

Notice that when I have a pointer to any base subobject of `Cat`, as long as that base class has a vptr, the vptr **will** point to a vtable of `Cat`.

So "what is the most derived type of this object" is answered by looking at the vtable.

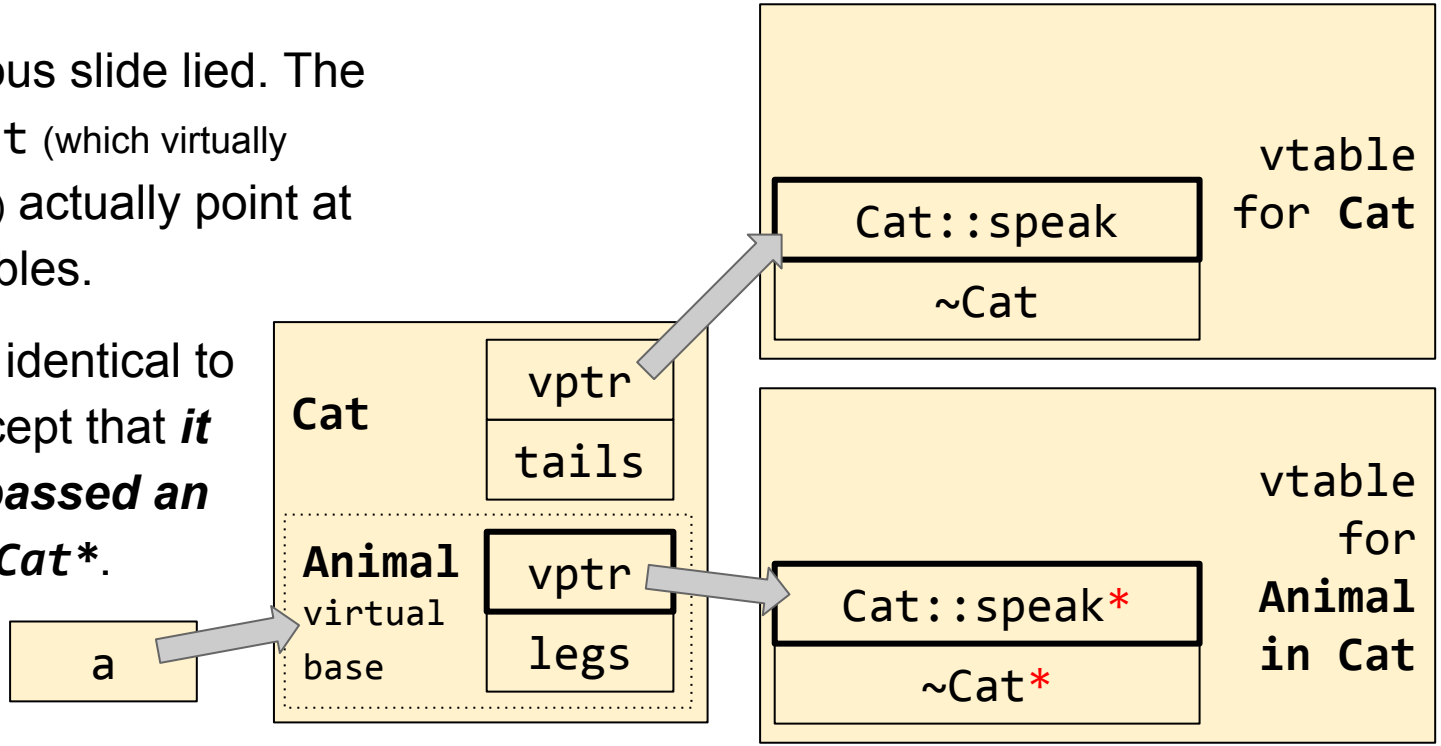


Secondary vtables and thunks

Okay, our previous slide lied. The two vptrs in a `Cat` (which virtually inherits from `Animal`) actually point at two different vtables.

`Cat::speak*` is identical to `Cat::speak` except that ***it expects to be passed an `Animal*`, not a `Cat*`***.

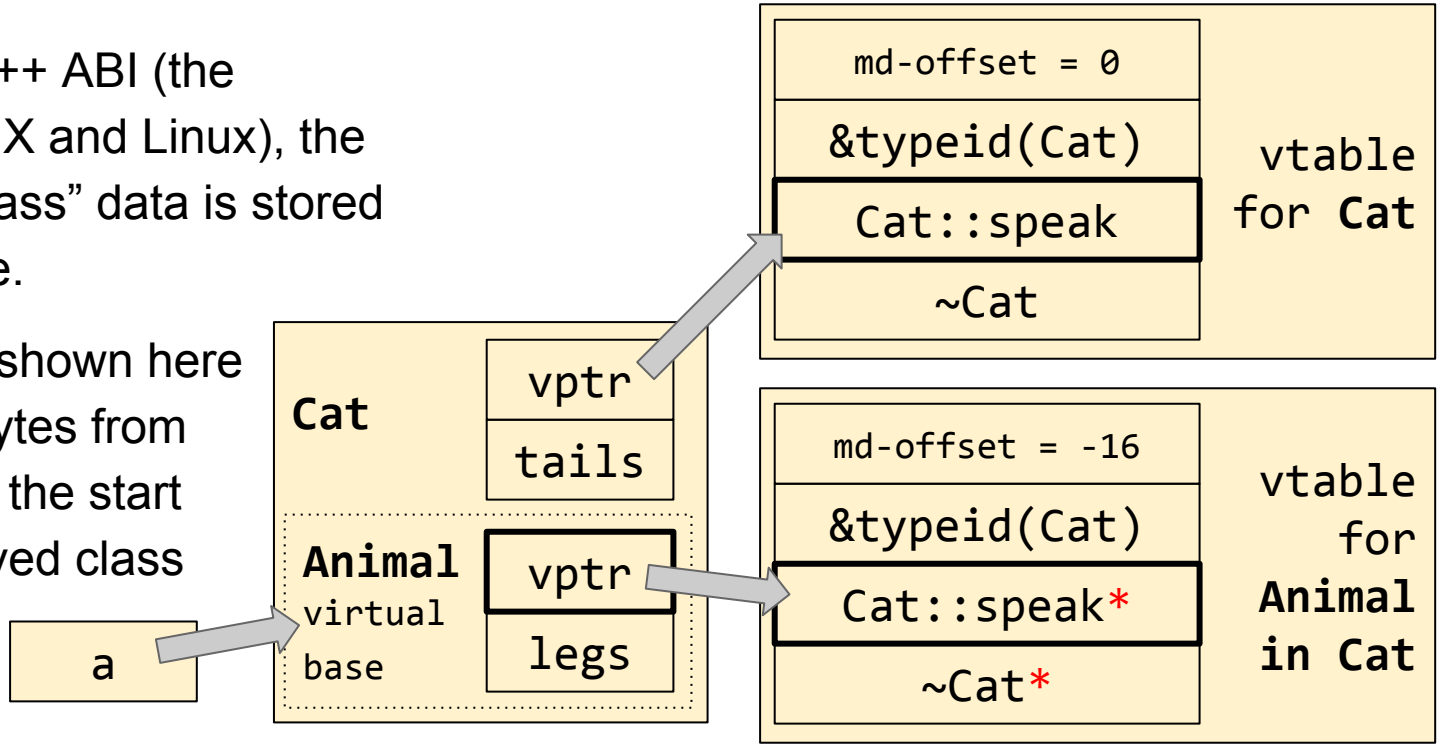
```
a->speak();
```



When in doubt: it's in the vtable

In the Itanium C++ ABI (the standard on OS X and Linux), the "most derived class" data is stored *before* the vtable.

The "md-offset" shown here is the offset in bytes from the given vptr to the start of the most derived class — in this case, `Cat`.

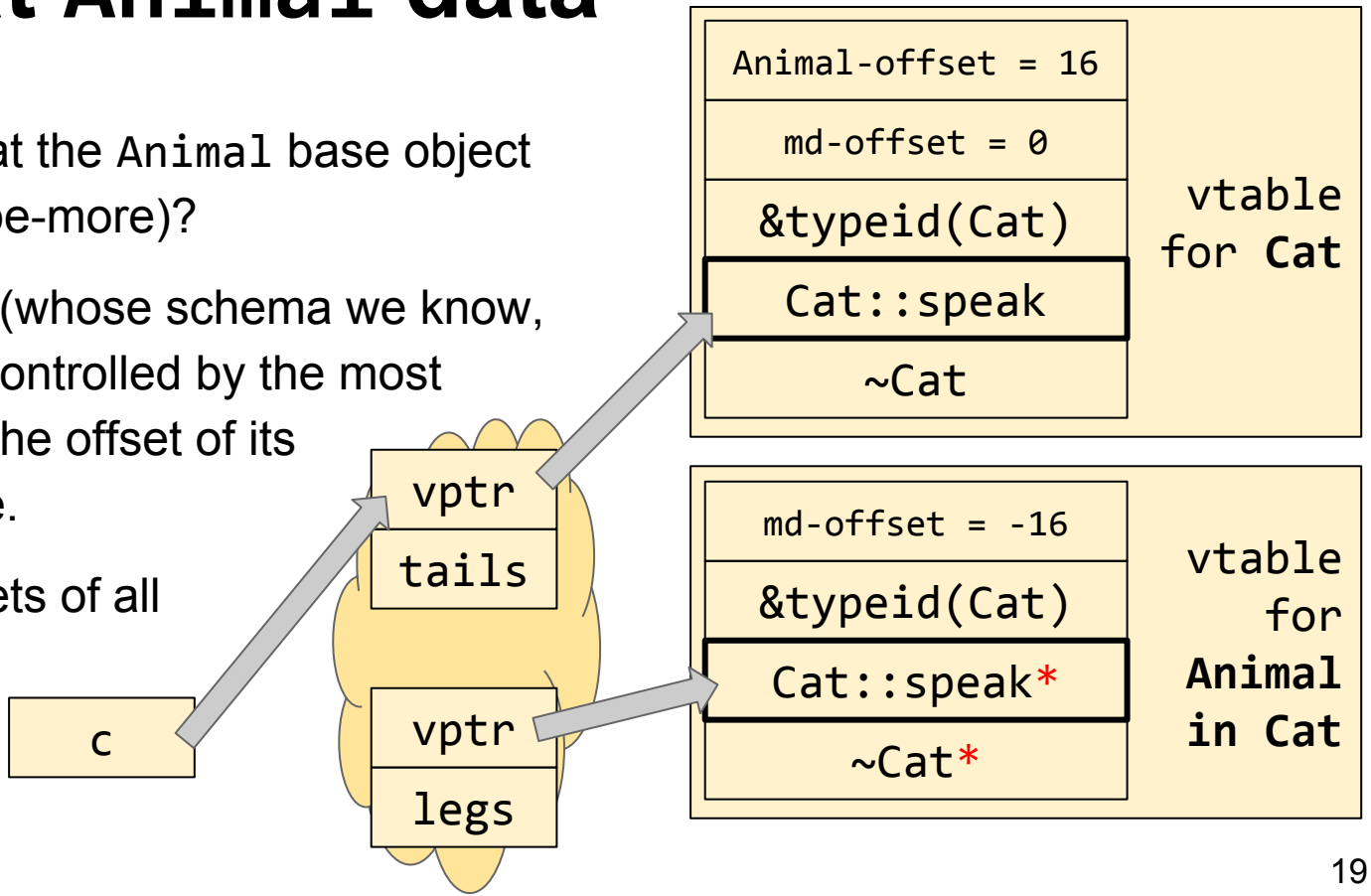


Getting at `Animal` data

So how do we get at the `Animal` base object of a `Cat` (and-maybe-more)?

Ask our `Cat` vtable (whose schema we know, but whose data is controlled by the most derived object) for the offset of its `Animal` virtual base.

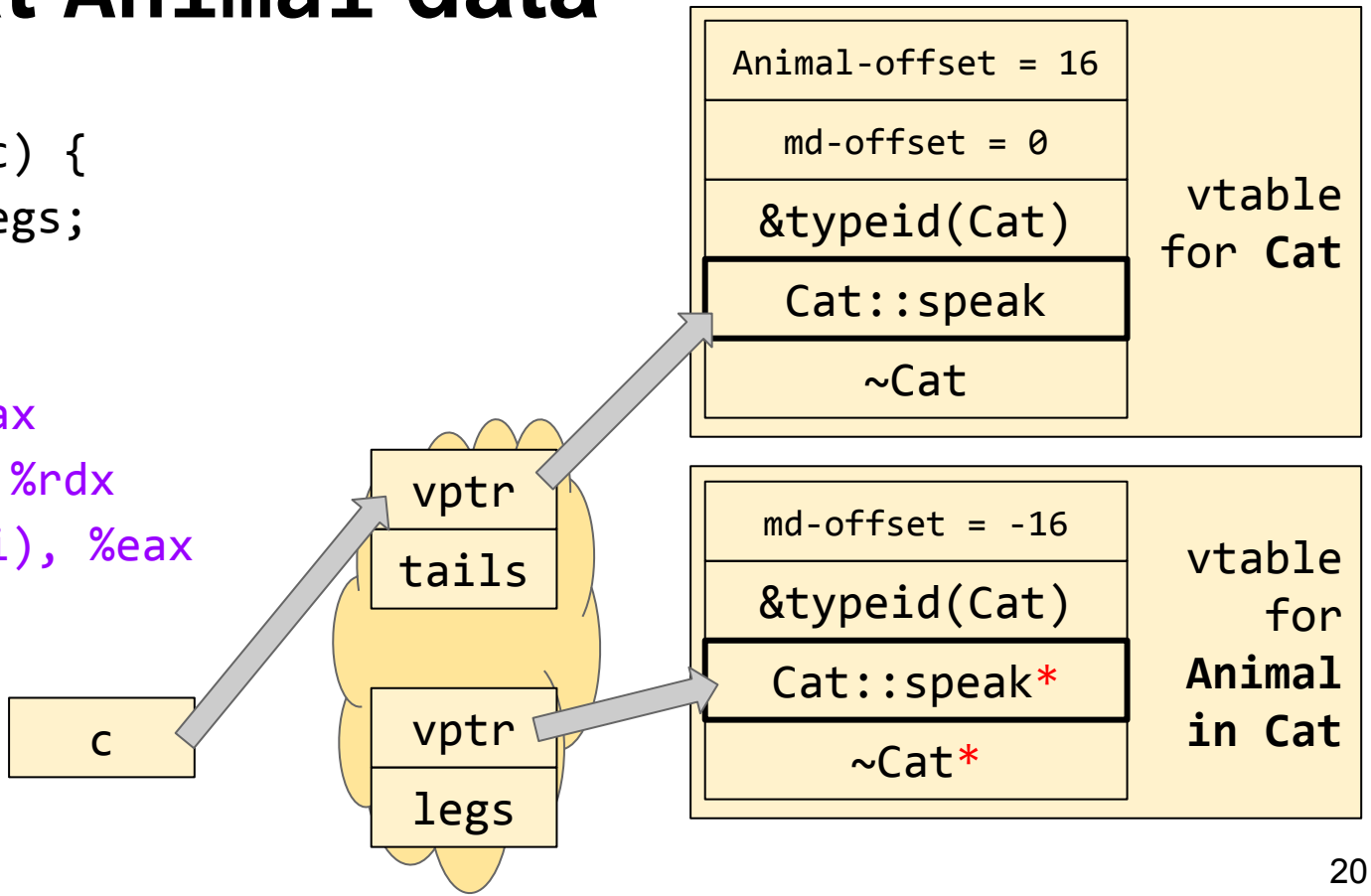
We'll store the offsets of all our virtual bases ahead of the md-offset.



Getting at `Animal` data

```
void test(Cat *c) {
    return c->legs;
}
```

```
movq (%rdi), %rax
movq -24(%rax), %rdx
movl 8(%rdx,%rdi), %eax
```



Vtable layout recap (Itanium ABI)

...

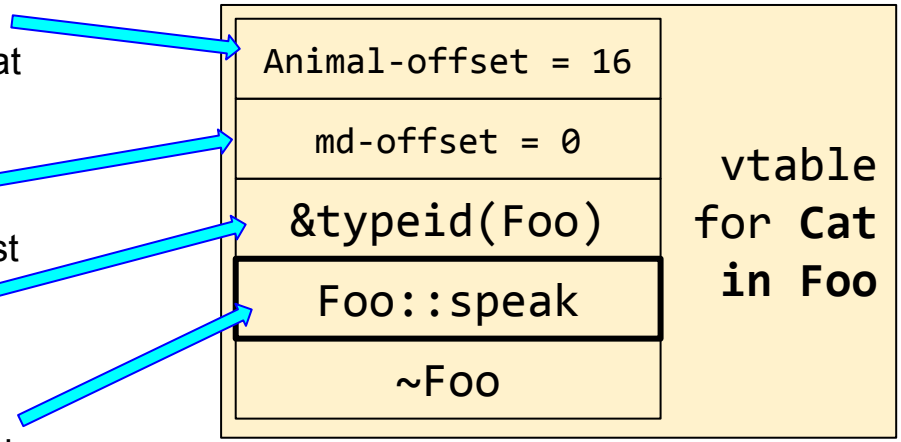
offsets (within the most derived object) to virtual bases of Cat

offset to the most derived object

type_info for the most derived object

pointers to the most derived object's versions of the virtual methods declared by Cat

...



`dynamic_cast` from scratch

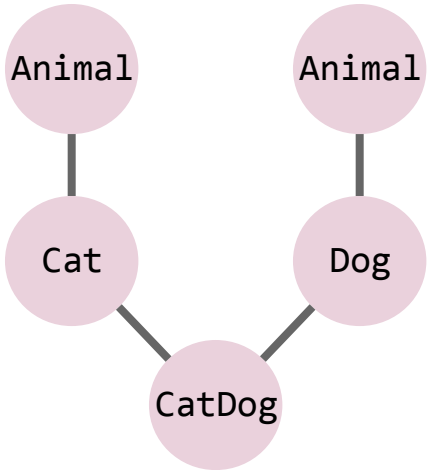
Two kinds of inheritance graphs

How I visualize complex class layout

```
struct Animal { virtual ~Animal(); }
struct Cat : public Animal {};
struct Dog : public Animal {};
struct CatDog : public Cat, Dog {};
```

Lines indicate inheritance.

Each “node” indicates a subobject.



How I visualize complex class layout

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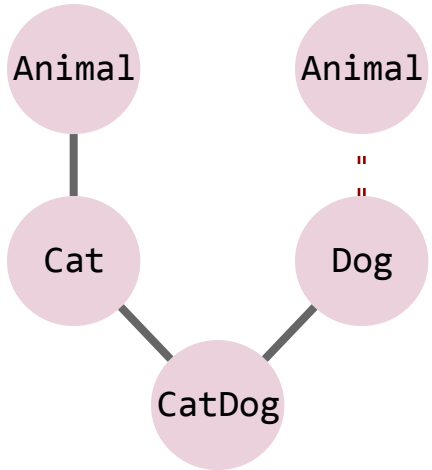

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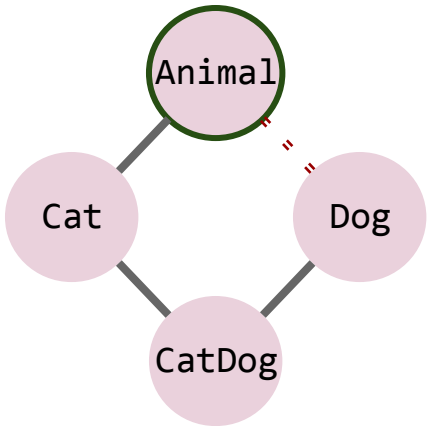

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Heavy circles indicate virtual inheritance.

- Notice that there is only ever *one* virtual subobject with a given name.

How I visualize complex class layout

Lines indicate inheritance.

Each “node” indicates a subobject.

Dotted lines indicate non-public inheritance.

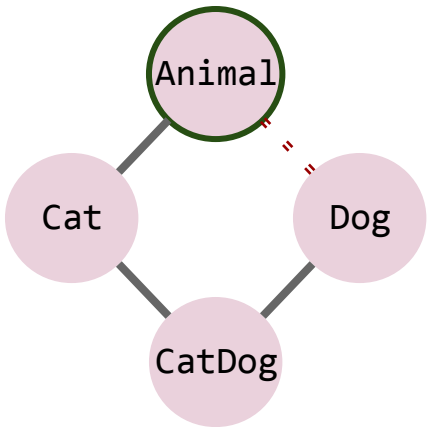
Heavy circles indicate virtual inheritance.

- Notice that there is only ever *one* virtual subobject with a given name.

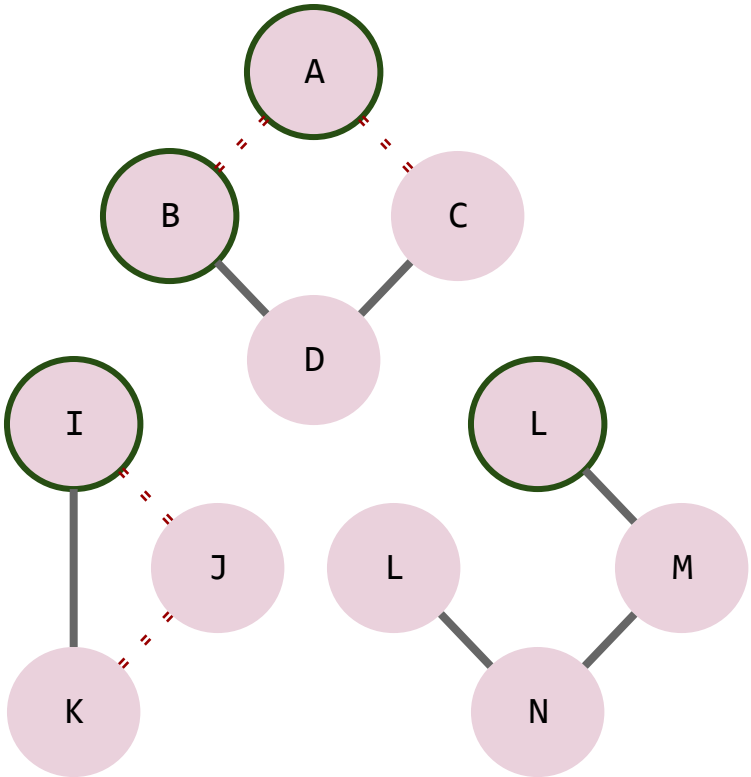
The root of the graph is the *most derived object*.

- Notice that there is always a single root.

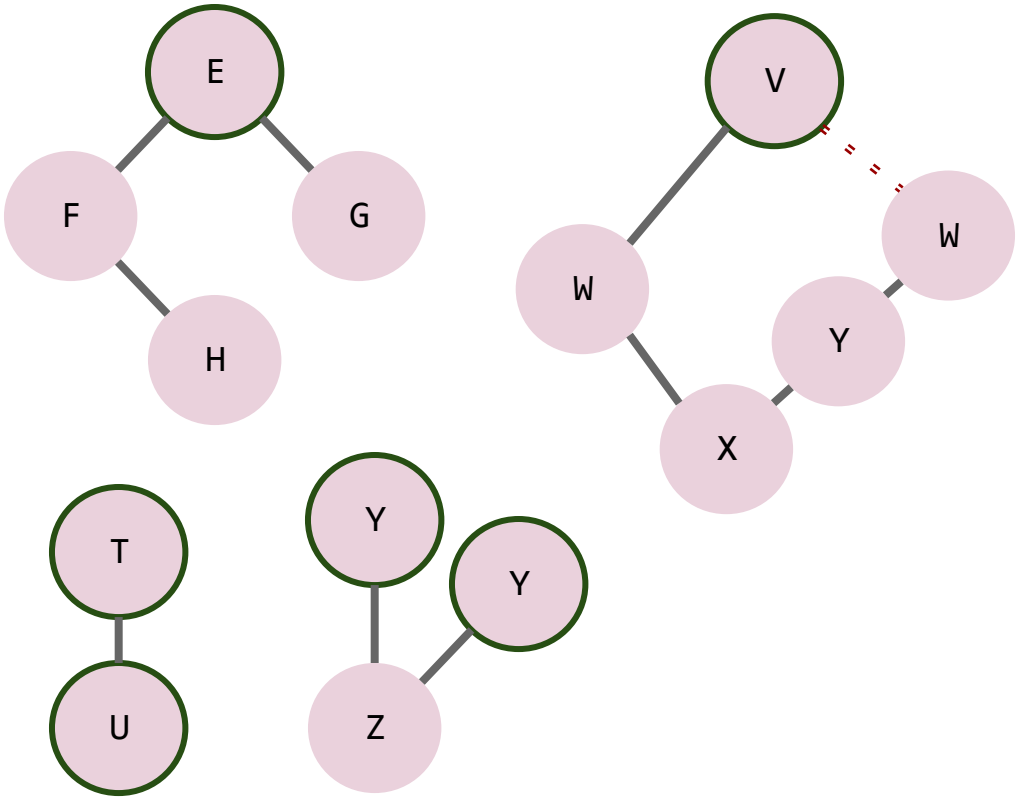
We don’t have a great way to indicate inheritance order, but that’s okay, because inheritance order doesn’t matter to (a correct implementation of) `dynamic_cast`.



Possible graphs



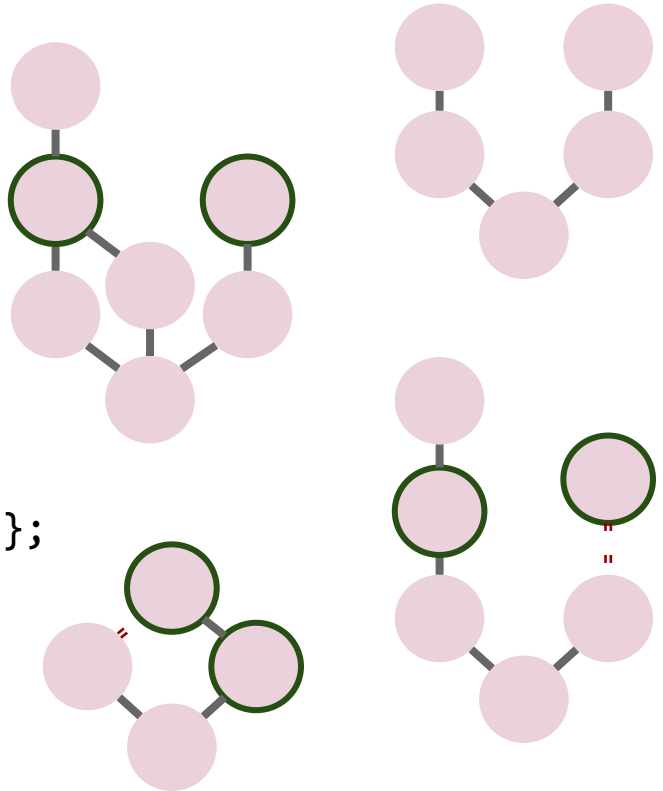
Impossible graphs



Match these graphs to their code

```
struct Animal;
struct Cat : Animal {};
struct Dog : Animal {};
struct Sponge : protected virtual Animal {};
struct LeftCat : virtual Cat {};
struct RightCat : virtual Cat {};
struct Flea : virtual Animal {};

struct CatDog : Cat, Dog {};
struct SiameseCat : LeftCat, RightCat, Flea {};
struct Bath : LeftCat, Sponge {};
struct Nemo : Sponge, virtual Flea {};
```



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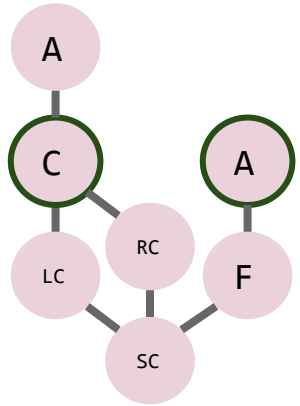

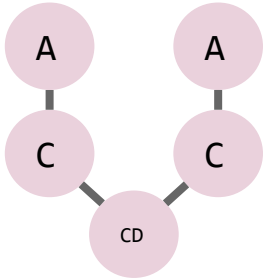

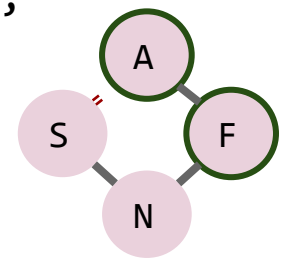

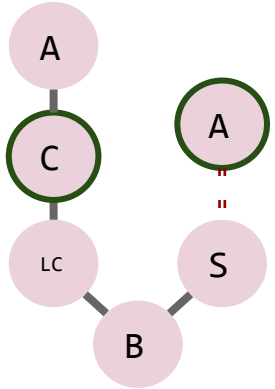

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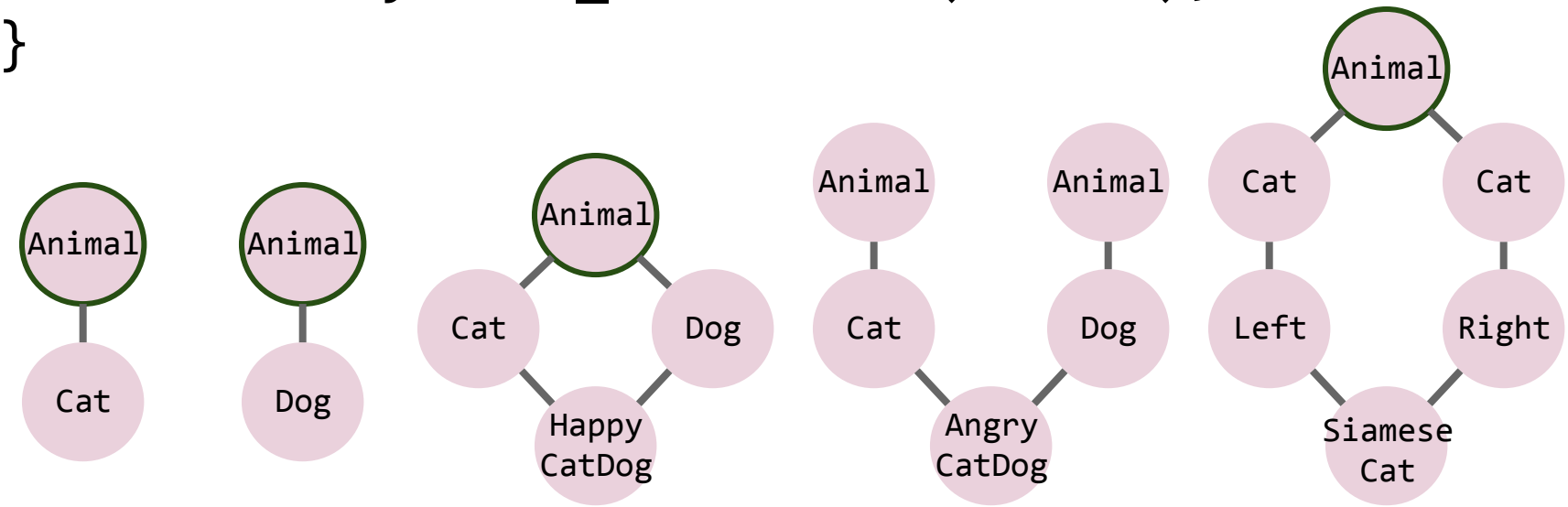
  - How to visualize complex class layouts [22–29]
- **Part II: `dynamic_cast`**
  - What should `dynamic_cast` do? [31–54]
  - Okay, how do we implement that? [55–60]
  - Benchmark numbers [61–64]

What should dynamic_cast do?

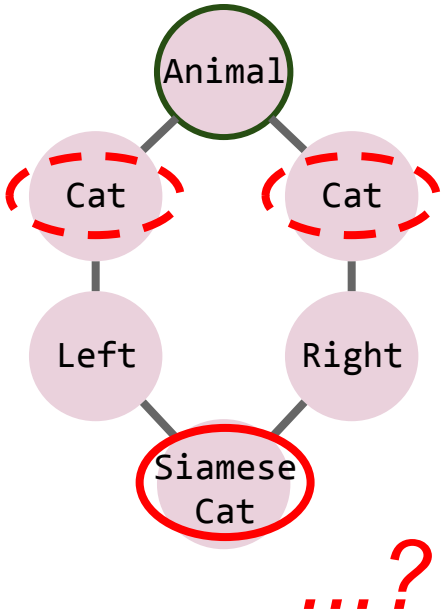
```
auto test(Animal *animal) {
    return dynamic_cast<Cat*>(animal);
}
```



What should `catch` do?

Ambiguous bases are basically invisible to RTTI.
This also applies during catch-handler matching.

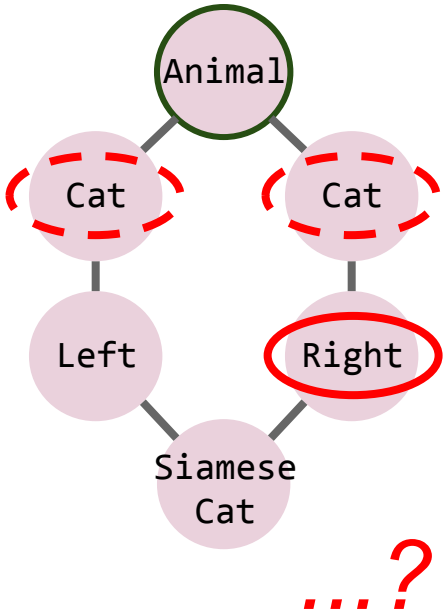
```
int main() {
    try {
        throw SiameseCat();
    } catch (const Cat&) {
        puts("SiameseCat IS-NOT-A Cat...");
        // it's two cats!
    } catch (const Animal&) {
        puts("...but SiameseCat IS-AN Animal!");
    }
}
```



What should this cast do?

This raises one more question: What about this case?

```
Cat *derived_to_sibling_or_to_base(RightCat *a) {
    return dynamic_cast<Cat *>(a);
}
```

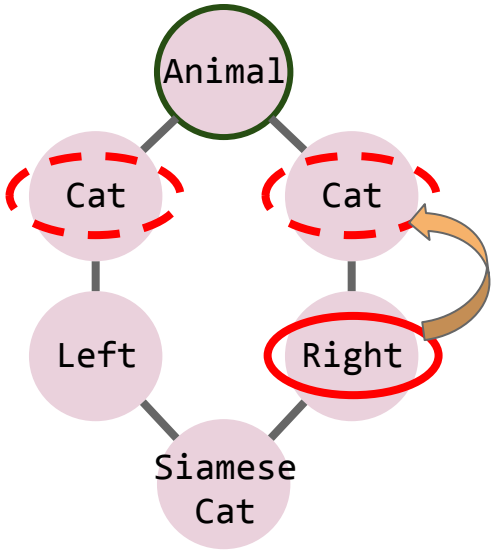


What should this cast do?

```
Cat *derived_to_sibling_or_to_base(RightCat *a)
{
    return dynamic_cast<Cat *>(a);
}
```

Trick question, sort of. This is not a “truly dynamic cast.” Because `Cat` is a known base class of `RightCat`, this cast is 100% equivalent to a `static_cast`. The compiler verifies that `Cat` is an accessible, unambiguous base class of `RightCat` and then generates the optimal code.

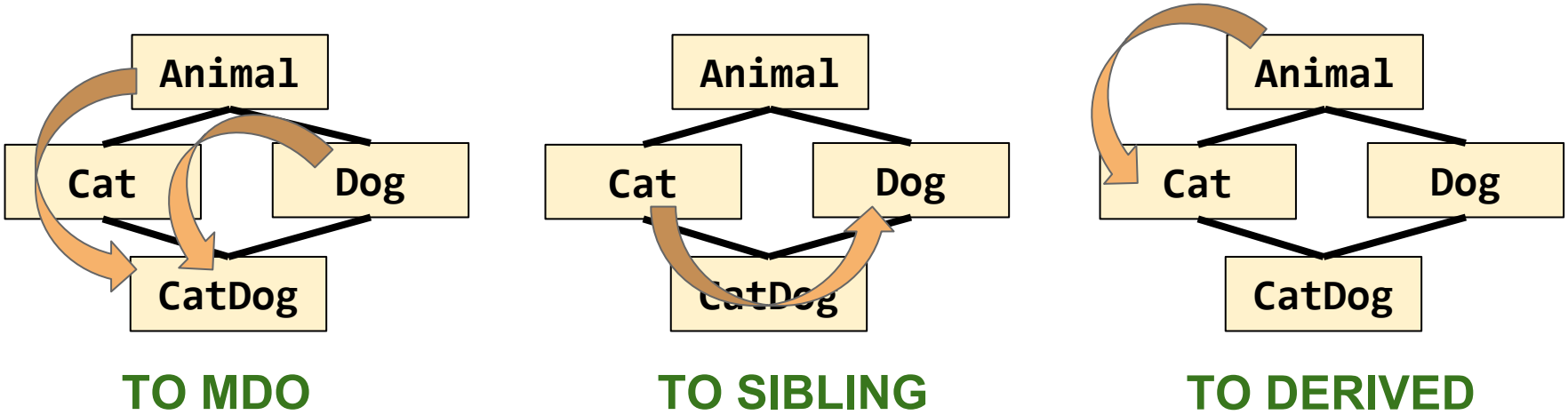
If `Cat` is an *inaccessible* base class of `RightCat`, the dynamic_cast is ill-formed and you get a compile-time error.



`dynamic_cast`, take 1

Per the Itanium spec, there are just three “truly dynamic casts” —

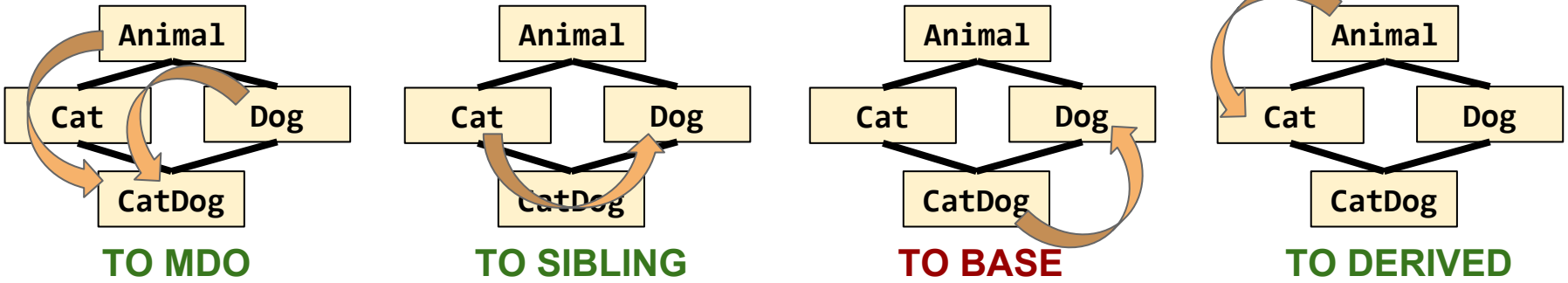
- `dynamic_cast<void*>` to the most-derived class
- `dynamic_cast` across the hierarchy, to a sibling base
- `dynamic_cast` from base to derived



The fourth “truly dynamic cast”

We need this one when deciding whether to stop stack-unwinding at a given `catch` clause. Let’s call it `castToBase`.

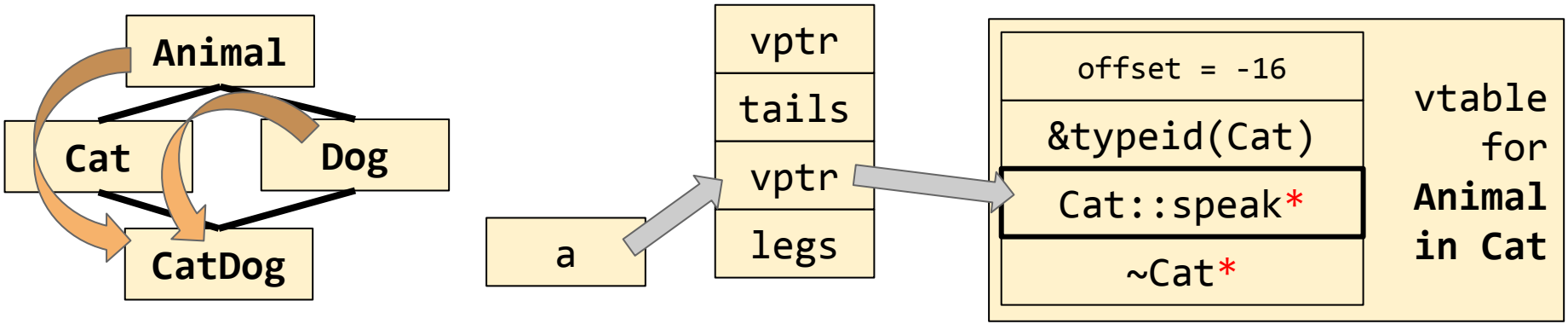
- `dynamic_cast<void*>` to the most-derived class
- `dynamic_cast` across the hierarchy, to a sibling base
- *The dynamic MDO-to-public-base conversion implied by `catch(Base&)`*
- `dynamic_cast` from base to derived



`dynamic_cast` to most-derived class

`dynamic_cast<void*>(p)` is literally just "load the offset-to-most-derived word located at `vptr[-2]` and add it to the `this` pointer."

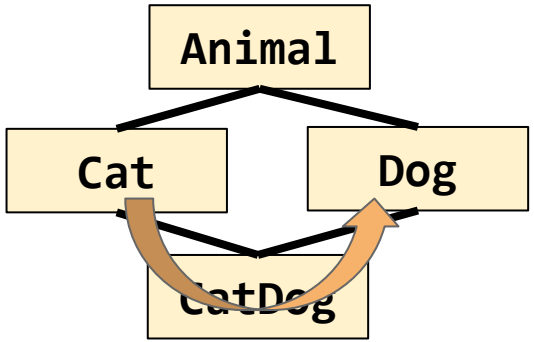
Every major compiler will generate optimal inline code for this.



DONE. SO EASY

`dynamic_cast` to sibling base

In this case we have a `Cat&` and we're trying to convert it to a `Dog&`.
The compiler already knows, statically, that `Cat` and `Dog` are "unrelated."
That is, not all `Cats` are `Dogs` and not all `Dogs` are `Cats`.
So the *only* way for this `dynamic_cast` to succeed is if somebody has made a `CatDog`.

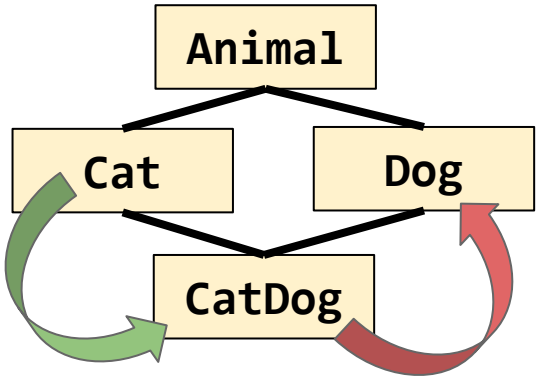


Clang optimizes iff `Cat` is final.
ICC, GCC, and MSVC do not optimize.

TO SIBLING

`dynamic_cast` to sibling base

- First, get a pointer to the most-derived class.
- If that most-derived class is still `Cat`, we can fail.
- Otherwise, "`castToBase`"!

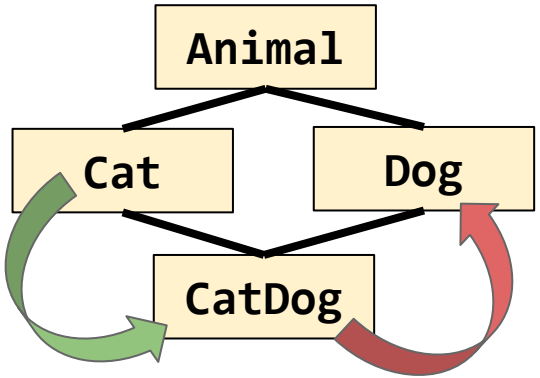


CAST TO MDO + CAST TO BASE

`dynamic_cast` MDO to base

We can imagine putting a compiler-generated function at a well-known offset in every vtable (or putting such a function pointer into every `std::type_info`):

```
void *castToBase(void *mdo, const type_info& to);
```



CAST TO MDO + CAST TO BASE

dynamic_cast MDO to base

```
void *Animal_castToBase(char *animal, const type_info& to) {
    return nullptr;
}
void *Cat_castToBase(char *cat, const type_info& to) {
    if (to == typeid(Animal)) return cat + 16;
    return nullptr;
}
void *HappyCatDog_castToBase(char *catdog, const type_info& to) {
    if (to == typeid(Cat)) return catdog + 0;
    if (to == typeid(Dog)) return catdog + 16;
    if (to == typeid(Animal)) return catdog + 32;
    return nullptr;
}
```

dynamic_cast MDO to base

```
void *Animal_castToBase(char *animal, const type_info& to) {
    return nullptr;
}
void *Cat_castToBase(char *cat, const type_info& to) {
    if (to == typeid(Animal)) return cat + 0;
    return nullptr;
}
void *AngryCatDog_castToBase(char *catdog, const type_info& to) {
    if (to == typeid(Cat)) return catdog + 0;
    if (to == typeid(Dog)) return catdog + 24;
    return nullptr;
}
```

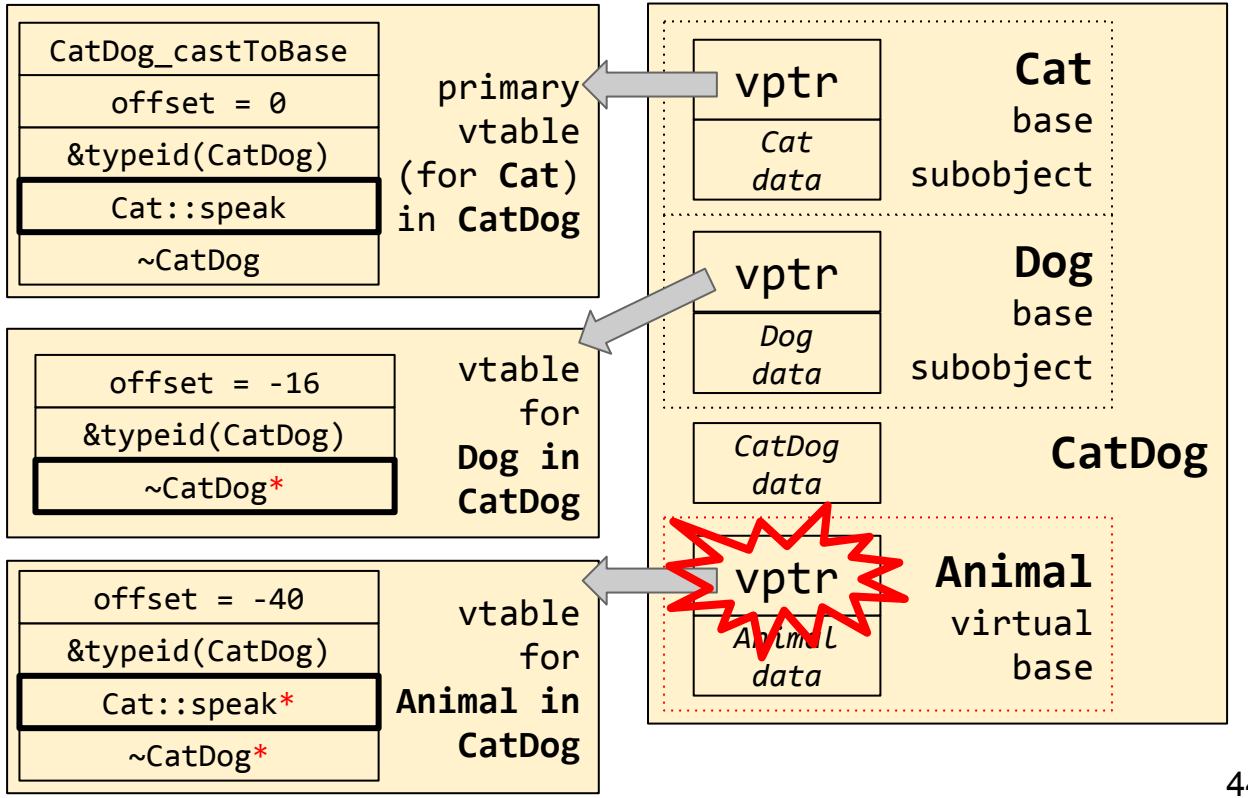
Let’s do an example.

Casting `Animal` to `Dog` (in a `CatDog`)

dynamic_cast<Dog*>(Animal*)

```
Dog *test(Animal *p) {
    return dynamic_cast<
        Dog*>(p);
}
```

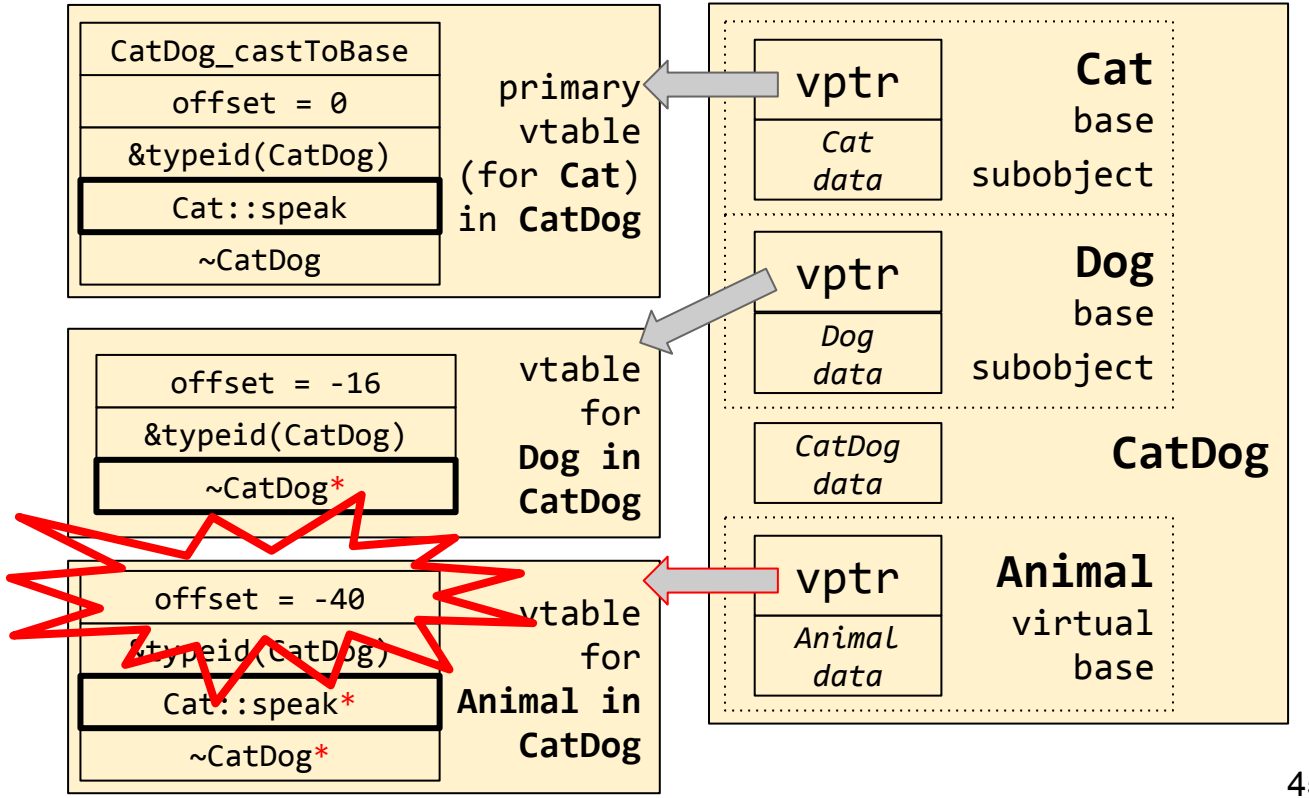
1. Get the vptr from p.



dynamic_cast<Dog*>(Animal*)

```
Dog *test(Animal *p) {
    return dynamic_cast<
        Dog*>(p);
}
```

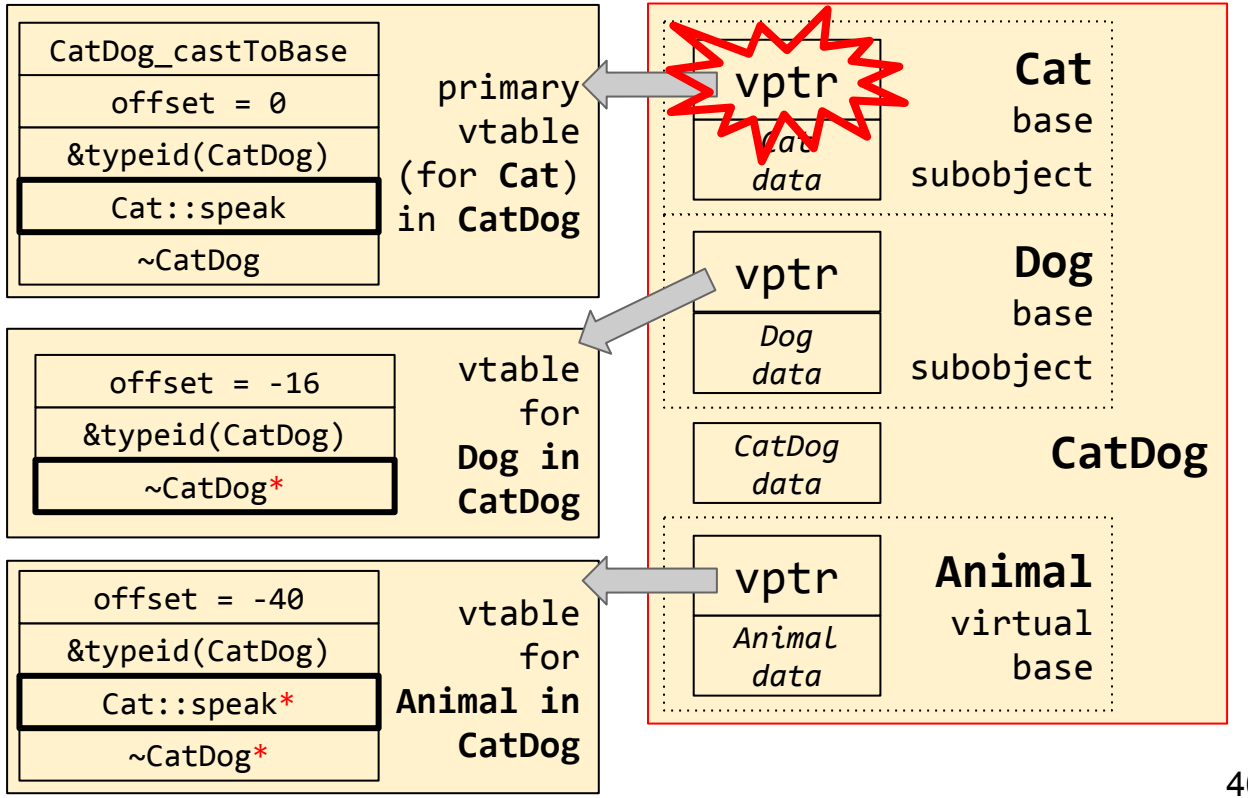
1. Get the vptr from `p`.
2. Load the offset-to-most-derived from `vptr[-2]` and adjust the `this` pointer.



dynamic_cast<Dog*>(Animal*)

```
Dog *test(Animal *p) {
    return dynamic_cast<
        Dog*>(p);
}
```

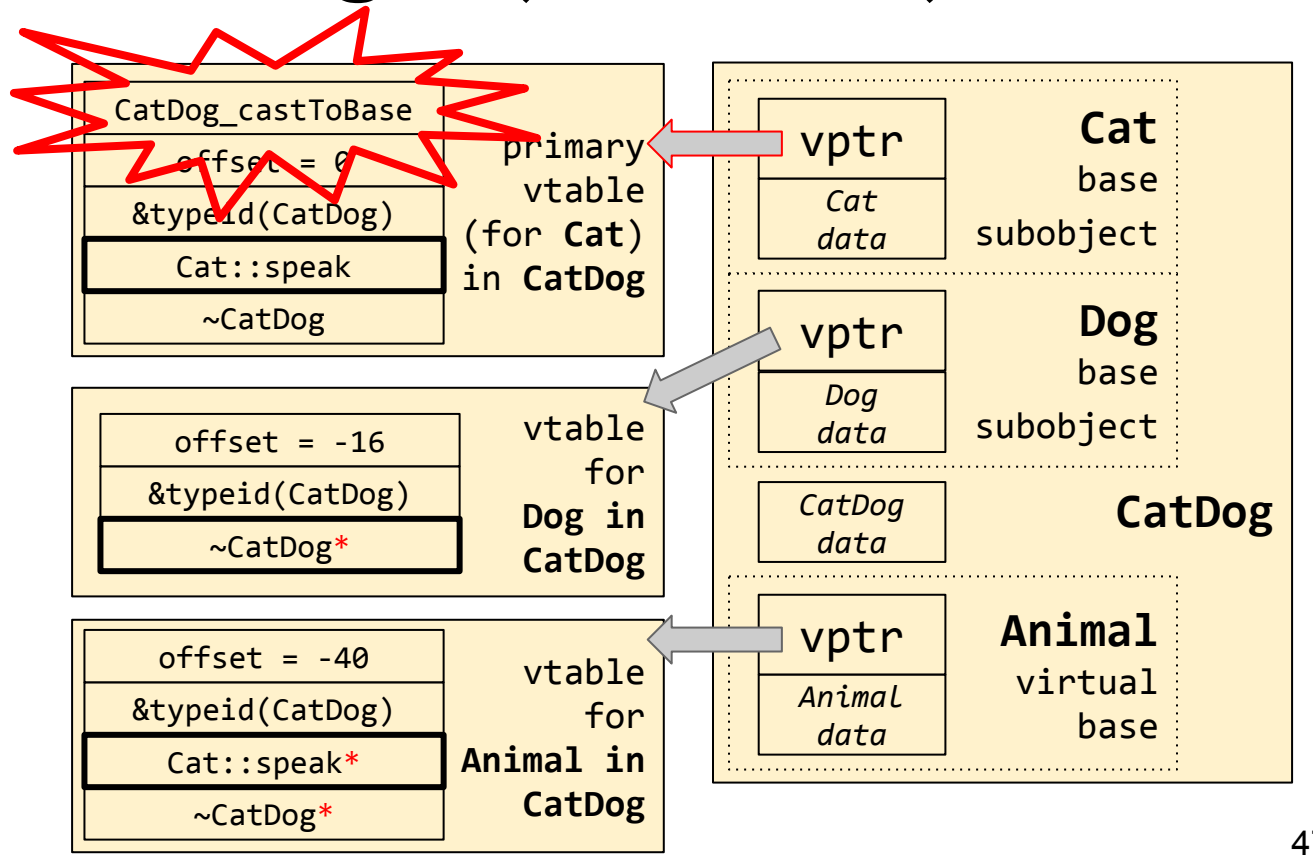
1. Get the vptr from `p`.
2. Load the offset-to-most-derived from `vptr[-2]` and adjust the `this` pointer.
3. Load the new vptr.



dynamic_cast<Dog*>(Animal*)

```
Dog *test(Animal *p) {
    return dynamic_cast<
        Dog*>(p);
}
```

1. Get the vptr from `p`.
2. Load the offset-to-most-derived from `vptr[-2]` and adjust the `this` pointer.
3. Load the new vptr.
4. Load the address of `castToBase` from `vptr[-3]`.
5. Call that function with `adjusted_this`, `typeid(Dog)`.



Wind up in `CatDog_castToBase`

```
void *CatDog_castToBase(void *p, const type_info& to) {
    if (to == typeid(Cat)) return p + 0;
    if (to == typeid(Dog)) return p + 16;
    if (to == typeid(Animal)) return p + 32;
    return nullptr;
}
```

This concludes the cases “truly dynamic cast to public base” and “cast between unrelated sibling bases.”

But actually this is NOT how either Itanium or MSVC do it! They encode the entire class hierarchy in data and run expensive code to parse that class hierarchy.
Why? Historical reasons, I guess.

“Accessible” versus “public”

```
struct Animal { virtual ~Animal(); };

struct Sponge : protected Animal {
    auto accessible(Sponge *s) {
        return static_cast<Animal*>(s);  // okay
    }
};
auto inaccessible(Sponge *s) {
    return static_cast<Animal*>(s);  // error
}
```

Either `static_cast` or `dynamic_cast` from derived to base requires that the target base be unambiguous and *accessible* in the current lexical scope.

“Accessible” versus “public”

```
struct Animal { virtual ~Animal(); };

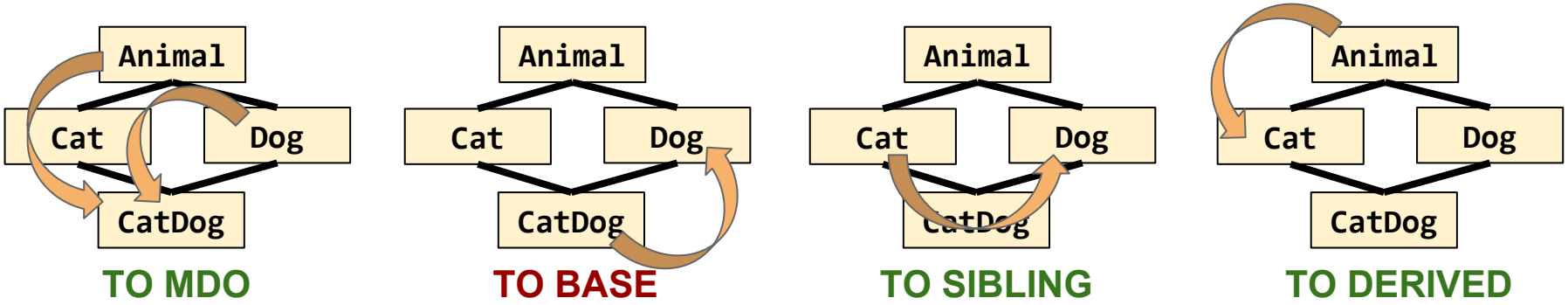
struct Sponge : protected Animal {
    void accessible(Sponge *s) {
        try { throw s; } catch (Animal *a) {}  // doesn't catch
    }
};
auto inaccessible(Sponge *s) {
    try { throw s; } catch (Animal *a) {}  // doesn't catch
}
```

`castToBase` (used by `catch` and `dynamic_cast` to sibling base) requires that the target base be unambiguous and *public* — there’s no such thing as lexical scope at runtime!

`dynamic_cast` from parent to child

- `dynamic_cast<void*>` to the most-derived class
- The dynamic derived-to-base conversion implied by `catch(Base&)`
- `dynamic_cast` across the hierarchy, to a sibling base
- `dynamic_cast` *from base to derived*

Casting from base to derived cannot use the same “cast to most-derived, then `castToBase`” trick that we used for sibling bases.

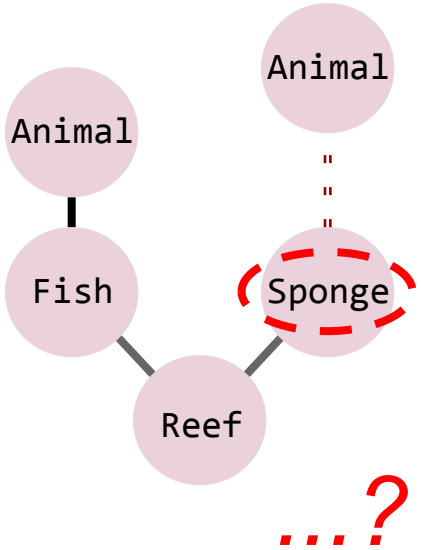


`dynamic_cast` from parent to child

Casting from base to derived cannot use the same “cast to most-derived, then `castToBase`” trick that we used for sibling bases. Consider:

```
struct Sponge : protected Animal {};
struct Fish : public Animal {};
struct Reef : Fish, Sponge {};

Sponge *foo(Animal *animal) {
    return dynamic_cast<Sponge*>(animal);
}
```

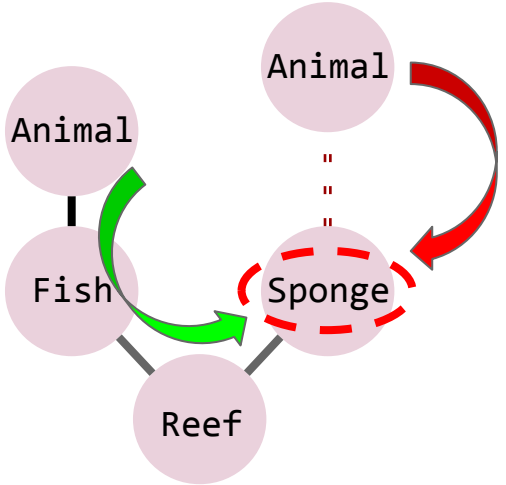


dynamic_cast from parent to child

Casting from base to derived cannot use the same “cast to most-derived, then castToBase” trick that we used for sibling bases. Consider:

```
struct Sponge : protected Animal {};
struct Fish : public Animal {};
struct Reef : Fish, Sponge {};

Sponge *foo(Animal *animal) {
    return dynamic_cast<Sponge*>(animal);
}
```



Our Sponge **IS-AN** Animal in certain lexical scopes, but at runtime there are no lexical scopes. So it’s *not* an Animal, and so, likewise, that Animal subobject doesn’t know it’s part of a Sponge.

`dynamic_cast` from parent to child

Start by finding the most derived object. (Get its RTTI from the primary vtable.) Ask it two questions:

- Your `Animal` subobject at offset *x* — is it a public base of a `Sponge` subobject? (If so, success: return the address of that `Sponge` subobject.)
- Your `Animal` subobject at offset *x* — is it a public base of yourself? (If so, delegate to `castToBase(mdo, typeid(Sponge))`.)

```
void *Reef_maybeFromHasAPublicChildOfTypeTo(void *mdo, int offset,
    const type_info& from, const type_info& to)
{
    if (offset == 0 && to == typeid(Fish)) return p + 0;
    // if (offset == 8 && to == typeid(Fish)) return nullptr;
    // if (offset == 8 && to == typeid(Sponge)) return nullptr;
    return nullptr;
}
```

Possible implementation

Again, the Itanium C++ ABI and the MSVC ABI take approaches that aren't this one. They use a graph of the class hierarchy, including "public/non-public" coloring of each edge, and they depth-first-search this graph on every `dynamic_cast`.

(And they do it wrong, with lots of bugs. Graph algorithms are hard.)

But let's put everything together and see what *we've* ended up with...

Possible implementation

```
template<class P, class From, class To = remove_pointer_t<P>>
To *dynamicast(From *p) {
    if constexpr (is_same_v<From, To>) {
        return p;
    } else if constexpr (is_base_of_v<To, From>) {
        return (To *)(p);
    } else if constexpr (is_void_v<To>) {
        return truly_dynamic_to_mdo(p);
    } else if constexpr (is_base_of_v<From, To>) {
        return truly_dynamic_from_base_to_derived<To>(p);
    } else {
        return truly_dynamic_between_unrelated_classes<To>(p);
    }
}
```

Technically we should also verify that `To` is an *accessible* base of `From` in the current lexical scope...

```
void *truly_dynamic_to_mdo(void *p)
{
    uint64_t *vptr =
        *reinterpret_cast<uint64_t **>(p);
    uint64_t mdoffset = vptr[-2];
    void *adjusted_this =
        static_cast<char *>(p) + mdoffset;
    return adjusted_this;
}

type_info& dynamic_typeid(void *p)
{
    return *(*reinterpret_cast<type_info ***>(p))[-1];
}
```

Fetch the vtable for "(some static type) in (most derived type)."

Get the "offset to MDO" from the vtable.

Get the MDO's typeinfo from offset -1 in the vtable.

```
template<class To, class From>
To *truly_dynamic_from_base_to_derived(From *p) {
    void *mdo = truly_dynamic_to_mdo(p);
    const type_info& ti = dynamic_typeid(mdo);
    int offset = (char *)p - (char *)mdo;
    if (ti == From) {
        return nullptr;
    } else if (ti == To && ti.isPublicBaseOfMe(offset, From)) {
        return (To *)(mdo);
    } else if (void *so = ti.maybeFromHasAPublicChildOfTypeTo(
                                  mdo, offset, From, To)) {
        return (To *)(so);
    } else if (ti.isPublicBaseOfMe(offset, From)) {
        return (To *)(ti.castToBase(mdo, To));
    }
    return nullptr;
}
```

Here, ***To*** represents the compile-time constant `typeid(To)`.

```
template<class To, class From>
To *truly_dynamic_between_unrelated_classes(From *p) {
    if constexpr (is_final_v<To> || is_final_v<From>) {
        return nullptr;
    }
    void *mdo = truly_dynamic_to_mdo(p);
    const type_info& ti = dynamic_typeid(mdo);
    int offset = (char *)p - (char *)mdo;
    if (ti == From) {
        return nullptr;
    } else if (ti == To) {
        return (To *)(mdo);
    } else if (ti.isPublicBaseOfMe(offset, From)) {
        return (To *)(ti.castToBase(mdo, To));
    }
    return nullptr;
}
```

Here, ***To*** represents the compile-time constant `typeid(To)`.

```
template<class P, class From, class To = remove_pointer_t<P>>
To *dynamicast(From *p) {
    if constexpr (is_same_v<From, To>) {
        return p;
    } else if constexpr (is_base_of_v<To, From>) {
        return (To *)(p);
    } else if constexpr (is_void_v<To>) {
        return truly_dynamic_to_mdo(p);
    } else if constexpr (is_base_of_v<From, To>) {
        return truly_dynamic_from_base_to_derived<To>(p);
    } else {
        return truly_dynamic_between_unrelated_classes<To>(p);
    }
}
```

Benchmark numbers

I wrote a Python script to generate random class hierarchies of 10 classes (of depth up to 7). I used Google Benchmark to measure `dynamic_cast` versus `dynamicast` for all possible combinations of `From*` and `To*`
(if the cast would be well-formed, not necessarily if it would return non-null).

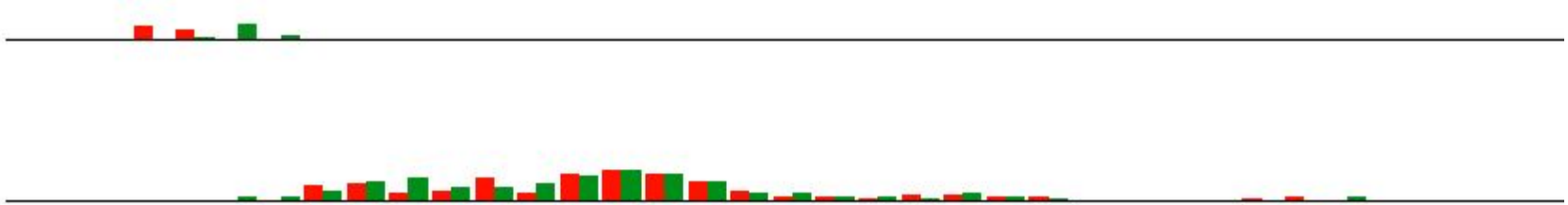
The code is on my GitHub: http://github.com/Quuxplusone/from-scratch/

```
template<class To, class From, bool_if_t<can_dynamic_cast_v<From, To>> = true>
void run_benchmark(From *f) {
    To *p = dynamic_cast<To*>(f);
    benchmark::DoNotOptimize(p);
}

template<class To, class From, bool_if_t<!can_dynamic_cast_v<From, To>> = true>
void run_benchmark(From *) {}
```

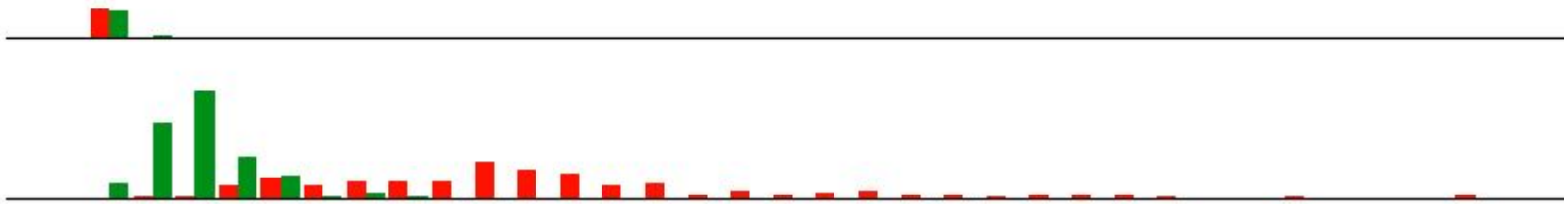
Benchmark numbers

Benchmark	Time	CPU	Iterations
BM_native_void	524 ns	522 ns	1312385
BM_dynamicast_void	973 ns	968 ns	731743
BM_native_Class1	1686 ns	1685 ns	393781
BM_dynamicast_Class1	1669 ns	1641 ns	400197
BM_native_Class2	2388 ns	2306 ns	287183
BM_dynamicast_Class2	2174 ns	2149 ns	307292
BM_native_Class3	2097 ns	2078 ns	339008
BM_dynamicast_Class3	1897 ns	1875 ns	363459



Benchmark numbers (-O3)

```
Benchmark                        Time             CPU Iterations
----------------------------------------------------------------
BM_native_void                  64 ns         64 ns    10421164
BM_dynamicast_void             124 ns        124 ns     5743071
BM_native_Class1              2303 ns       2296 ns      319906
BM_dynamicast_Class1           422 ns        420 ns     1608478
BM_native_Class2              1941 ns       1937 ns      376849
BM_dynamicast_Class2           416 ns        416 ns     1760156
BM_native_Class3              1125 ns       1124 ns      641425
BM_dynamicast_Class3           233 ns        232 ns     3067391
```



Benchmark numbers (-O3 only for typeinfo)

Benchmark	Time	CPU	Iterations
BM_native_void	459 ns	455 ns	1676245
BM_dynamicast_void	665 ns	663 ns	1107805
BM_native_Class1	2491 ns	2412 ns	311500
BM_dynamicast_Class1	1222 ns	1177 ns	640439
BM_native_Class2	2470 ns	2377 ns	301865
BM_dynamicast_Class2	1136 ns	1134 ns	664225
BM_native_Class3	1453 ns	1452 ns	491394
BM_dynamicast_Class3	793 ns	793 ns	922704