

# Introduction to Quantum Game Theory

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# Research Questions

## Mathematical description of Quantum Game Theory

What are the main mathematical concepts of quantum game theory?  
How are the theories (Game Theory and Quantum Theory) unified?

## Results for Quantum Games within different game classes

What are the main differences between classical and quantum game theory. Is the underlying Nash equilibrium structure of (2 player)-(2 strategy) games changed within a quantum game theory-based analysis?

## Presentation of various applications

How can quantum game theory be applied to real game situations?





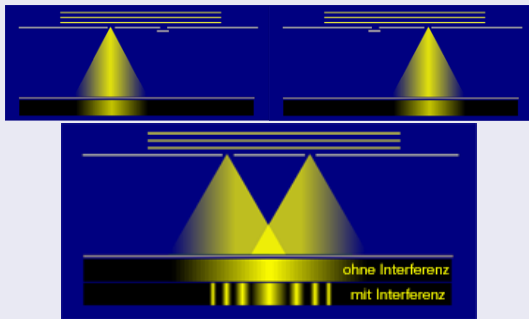






# Welle-Teilchen-Dualismus

## Das Doppelspaltexperiment



**Figure:** Beim Doppelspaltexperiment offenbaren Teilchen ihre Welleneigenschaften. Quelle: Michael Craiss

1961 wurde das Doppelspaltexperiment mit Elektronen durch Claus Jönsson durchgeführt und im September 2002 in einer Umfrage der englischen physikalischen Gesellschaft in der Zeitschrift 'Physics World' zum schönsten physikalischen Experiment aller Zeiten gewählt.







# Das Einstein-Podolsky-Rosen Paradoxon

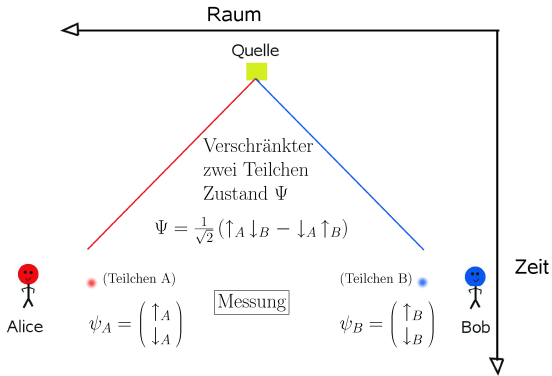


Figure: EPR Gedankenexperiment ('Bohm-Aharonov' Version). Der durch die Messung an Teilchen A verursachte instantane Kollaps der 2-Teilchen Wellenfunktion  $\Psi$  bestimmt den Zustand des Teilchens B.

# Die Quantenverschränkung

## Verschränkte 2-Niveau-Quantensysteme

Zwei Teilchen (A,B) haben die Möglichkeit zwischen zwei Zuständen ( $\uparrow$ ,  $\downarrow$ ) zu wählen. Die Basisvektoren der jeweiligen Hilberträume der Teilchen seien wie folgt definiert:

Zustand des Teilchens A:  $\psi_A \in \mathcal{H}_A \equiv \mathbb{C}^2$ , Basis:  $\{\uparrow_A, \downarrow_A\}$

Zustand des Teilchens B:  $\psi_B \in \mathcal{H}_B \equiv \mathbb{C}^2$ , Basis:  $\{\uparrow_B, \downarrow_B\}$

Der Hilbertraum des zusammengesetzten Systems ist ein komplexer vierdimensionaler Raum ( $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$ ). Der Gesamtzustand des 2-Teilchen Systems  $\Psi$  kann unter Umständen nicht in die jeweiligen Einzelzustände separiert werden.

Verschränkter Zustand (z.B.):  $\Psi = \frac{1}{\sqrt{2}} (\uparrow_A \downarrow_B - \downarrow_A \uparrow_B)$



# Definition of a (2 player)-(2 strategy) game $\Gamma$

An unsymmetric ( $2 \times 2$ ) game  $\Gamma$  is defined as ...

( $2 \times 2$ ) Game:  $\Gamma := (\{A, B\}, \mathcal{S}^A \times \mathcal{S}^B, \hat{\$}_A, \hat{\$}_B)$

Set of pure strategies of player A and B:  $\mathcal{S}^A = \{s_1^A, s_2^A\}, \mathcal{S}^B = \{s_1^B, s_2^B\}$

Set of mixed strategies of player A and B:  $\tilde{\mathcal{S}}^A = \{\tilde{s}_1^A, \tilde{s}_2^A\}, \tilde{\mathcal{S}}^B = \{\tilde{s}_1^B, \tilde{s}_2^B\}$

Payoff matrix for player A:  $\hat{\$}_A = \begin{pmatrix} \$_{11}^A & \$_{12}^A \\ \$_{21}^A & \$_{22}^A \end{pmatrix}$

Payoff matrix for player B:  $\hat{\$}_B = \begin{pmatrix} \$_{11}^B & \$_{12}^B \\ \$_{21}^B & \$_{22}^B \end{pmatrix}$



# The mixed strategy payoff function $\tilde{\$}^\mu$ of player $\mu = A, B$

Due to the normalizing conditions

$$\tilde{s}_1^\mu + \tilde{s}_2^\mu = 1 \quad \forall \mu = A, B$$

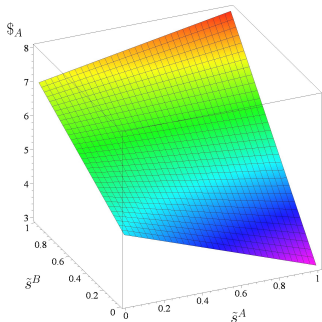
it is possible to simplify the functional dependence of the mixed strategy payoff function:

$$\begin{aligned}\tilde{\$}^\mu : ([0, 1] \times [0, 1]) &\rightarrow \mathbb{R} \\ \tilde{\$}^\mu(\tilde{s}^A, \tilde{s}^B) &= \$_{11}^\mu \tilde{s}^A \tilde{s}^B + \$_{12}^\mu \tilde{s}^A (1 - \tilde{s}^B) + \\ &\quad + \$_{21}^\mu (1 - \tilde{s}^A) \tilde{s}^B + \$_{22}^\mu (1 - \tilde{s}^A) (1 - \tilde{s}^B)\end{aligned}$$

, where  $\tilde{s}^A := \tilde{s}_1^A$ ,  $\tilde{s}^B := \tilde{s}_1^B$ ,  $\tilde{s}_2^A = 1 - \tilde{s}_1^A$  and  $\tilde{s}_2^B = 1 - \tilde{s}_1^B$

# The mixed strategy payoff function $\tilde{\$}^\mu$ of player $\mu = A, B$

Mixed strategy payoff function  $\tilde{\$}^A(\tilde{s}^A, \tilde{s}^B)$  of player A  
 ( $\$_{11}^A = 8, \$_{12}^A = 5, \$_{21}^A = 7, \$_{22}^A = 3$ )



Payoff  $\tilde{\$}^\mu(\tilde{s}^A, \tilde{s}^B)$  as a function of  $\tilde{s}^A, \tilde{s}^B \in [0, 1]$ :

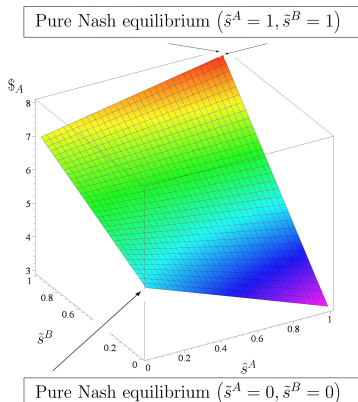
$$\begin{aligned} \tilde{\$}^\mu : ([0, 1] \times [0, 1]) &\rightarrow \mathbb{R} \\ \tilde{\$}^\mu(\tilde{s}^A, \tilde{s}^B) &= \$_{11}^\mu \tilde{s}^A \tilde{s}^B + \$_{12}^\mu \tilde{s}^A (1 - \tilde{s}^B) \\ &+ \$_{21}^\mu (1 - \tilde{s}^A) \tilde{s}^B + \$_{22}^\mu (1 - \tilde{s}^A)(1 - \tilde{s}^B) \\ &, \text{ where } \tilde{s}^A := \tilde{s}_1^A, \tilde{s}^B := \tilde{s}_1^B, \\ &\tilde{s}_2^A = 1 - \tilde{s}_1^A \text{ and } \tilde{s}_2^B = 1 - \tilde{s}_1^B \end{aligned}$$

Payoff  $\tilde{\$}^\mu(\tilde{\mathcal{S}}^A \times \tilde{\mathcal{S}}^B)$  as a function of the sets of mixed strategies for player A and B:

$$\begin{aligned} \tilde{\$}^\mu : (\tilde{\mathcal{S}}^A \times \tilde{\mathcal{S}}^B) &\rightarrow \mathbb{R} \\ \tilde{\$}^\mu((\tilde{s}_1^A, \tilde{s}_2^A), (\tilde{s}_1^B, \tilde{s}_2^B)) &= \$_{11}^\mu \tilde{s}_1^A \tilde{s}_1^B + \$_{12}^\mu \tilde{s}_1^A \tilde{s}_2^B \\ &+ \$_{21}^\mu \tilde{s}_2^A \tilde{s}_1^B + \$_{22}^\mu \tilde{s}_2^A \tilde{s}_2^B \end{aligned}$$

# Nash equilibria (NE)

## Nash equilibria and $\tilde{\$}^\mu(\tilde{s}^A, \tilde{s}^B)$



A strategy combination  $(\tilde{s}^{A*}, \tilde{s}^{B*})$  is called a Nash equilibrium, if:

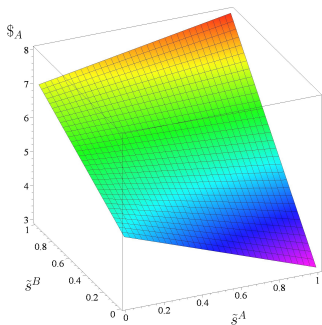
$$\begin{aligned}\tilde{\$}^A(\tilde{s}^{A*}, \tilde{s}^{B*}) &\geq \tilde{\$}^A(\tilde{s}^A, \tilde{s}^{B*}) \quad \forall \tilde{s}^A \in [0, 1] \\ \tilde{\$}^B(\tilde{s}^{A*}, \tilde{s}^{B*}) &\geq \tilde{\$}^B(\tilde{s}^{A*}, \tilde{s}^B) \quad \forall \tilde{s}^B \in [0, 1]\end{aligned}$$

A strategy combination  $(\tilde{s}^{A*}, \tilde{s}^{B*})$  is called an interior (mixed strategy) Nash equilibrium, if:

$$\begin{aligned}\left. \frac{\partial \tilde{\$}^A(\tilde{s}^A, \tilde{s}^B)}{\partial \tilde{s}^A} \right|_{\tilde{s}^B = \tilde{s}^{B*}} &= 0 \quad \forall \tilde{s}^A \in [0, 1], \tilde{s}^{B*} \in ]0, 1[ \\ \left. \frac{\partial \tilde{\$}^B(\tilde{s}^A, \tilde{s}^B)}{\partial \tilde{s}^B} \right|_{\tilde{s}^A = \tilde{s}^{A*}} &= 0 \quad \forall \tilde{s}^B \in [0, 1], \tilde{s}^{A*} \in ]0, 1[ \end{aligned}$$

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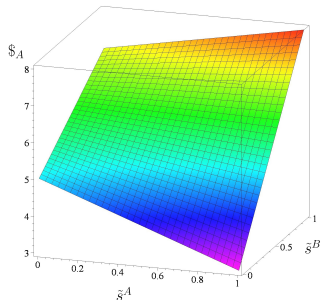
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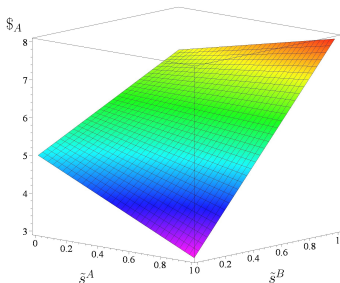
A strategy combination  $(\tilde{s}^{A*}, \tilde{s}^{B*})$  is called an interior (mixed strategy) Nash equilibrium, if:

$$\begin{aligned} \left. \frac{\partial \tilde{\$}^A(\tilde{s}^A, \tilde{s}^B)}{\partial \tilde{s}^A} \right|_{\tilde{s}^B = \tilde{s}^{B*}} &= 0 \quad \forall \tilde{s}^A \in [0, 1], \tilde{s}^{B*} \in ]0, 1[ \\ \left. \frac{\partial \tilde{\$}^B(\tilde{s}^A, \tilde{s}^B)}{\partial \tilde{s}^B} \right|_{\tilde{s}^A = \tilde{s}^{A*}} &= 0 \quad \forall \tilde{s}^B \in [0, 1], \tilde{s}^{A*} \in ]0, 1[ \end{aligned}$$



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# Replicator dynamics

Replicator dynamics: The dynamical behavior of a population of players

$$\frac{dx_i^A(t)}{dt} = x_i^A(t) \left[ \sum_{l=1}^2 \$_{il}^A x_l^B(t) - \sum_{l=1}^2 \sum_{k=1}^2 \$_{kl}^A x_k^A(t) x_l^B(t) \right]$$

$$\frac{dx_i^B(t)}{dt} = x_i^B(t) \left[ \sum_{l=1}^2 \$_{li}^B x_l^A(t) - \sum_{l=1}^2 \sum_{k=1}^2 \$_{lk}^B x_l^A(t) x_k^B(t) \right]$$

The two population vectors  $\vec{x}^A$  and  $\vec{x}^B$  have to fulfill the normalizing conditions of a unity vector

$$x_i^\mu(t) \geq 0 \quad \text{and} \quad \sum_{i=1}^2 x_i^\mu(t) = 1 \quad \forall i = 1, 2, t \in \mathbb{R}, \mu = A, B$$

# Replicator dynamics of $(2 \times 2)$ games

## Replicator dynamics of unsymmetric $(2 \times 2)$ games

$$\begin{aligned}\frac{dx(t)}{dt} &= \left( (\$_{11}^A + \$_{22}^A - \$_{12}^A - \$_{21}^A) (x(t) - (x(t))^2) \right) y(t) + (\$_{12}^A - \$_{22}^A) (x(t) - (x(t))^2) =: g_A(x, y) \\ \frac{dy(t)}{dt} &= \left( (\$_{11}^B + \$_{22}^B - \$_{12}^B - \$_{21}^B) (y(t) - (y(t))^2) \right) x(t) + (\$_{12}^B - \$_{22}^B) (y(t) - (y(t))^2) =: g_B(x, y)\end{aligned}$$

## Replicator dynamics of symmetric $(2 \times 2)$ games

$$\begin{aligned}\frac{dx}{dt} &= x \left[ \$_{11}(x - x^2) + \$_{12}(1 - 2x + x^2) + \$_{21}(x^2 - x) + \$_{22}(2x - x^2 - 1) \right] \\ &= x \left[ (\$_{11} - \$_{21})(x - x^2) + (\$_{12} - \$_{22})(1 - 2x + x^2) \right] =: g(x)\end{aligned}$$

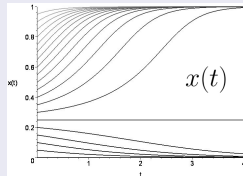
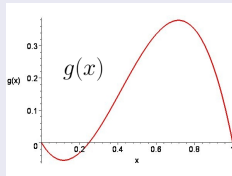
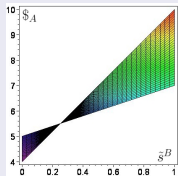
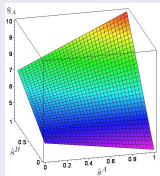
with:  $x = x(t) := x_1(t) \rightarrow x_2(t) = (1 - x(t))$



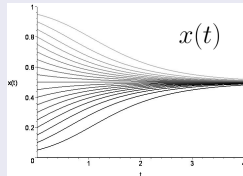
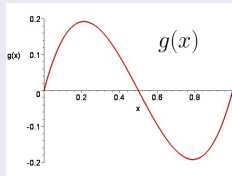
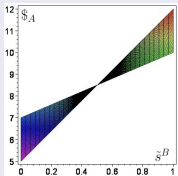
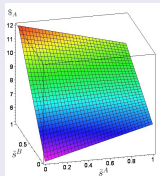


# Coordination $(a, b > 0)$ and Anti-Coordination $(a, b < 0)$ Class

Coordination game:  $a=3, b=1$ , two pure and one interior NE at  $\tilde{s}^* = \frac{1}{4}$ , two ESS  $((s_1^A, s_1^B)$  and  $(s_2^A, s_2^B))$



Anti-Coordination game:  $a=-2, b=-2$ , two pure asymmetric NE and one interior NE at  $\tilde{s}^* = \frac{1}{2}$ , one ESS  $(\tilde{s}^{A*} = \frac{1}{2}, \tilde{s}^{B*} = \frac{1}{2})$

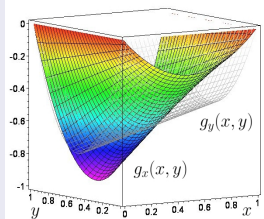




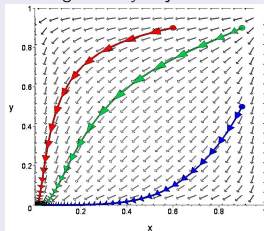
# Game classes of unsymmetric (2 player)-(2 strategy) games

## Corner Class (one ESS)

$g_x(x, y)$  (colored) and  $g_y(x, y)$  (wired):

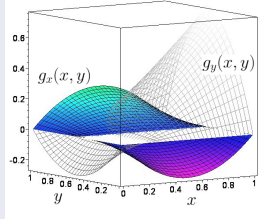


Phase diagram of  $xy$ -trajectories:

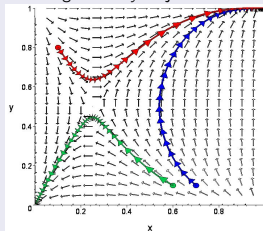


## Saddle Class (two ESS)

$g_x(x, y)$  (colored) and  $g_y(x, y)$  (wired):

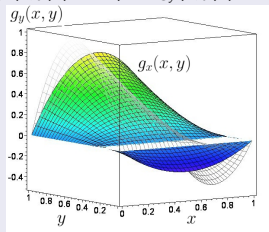


Phase diagram of  $xy$ -trajectories:

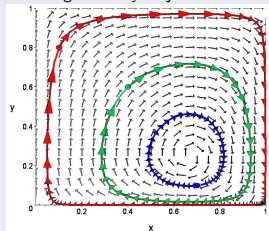


## Center Class (no ESS)

$g_x(x, y)$  (colored) and  $g_y(x, y)$  (wired):



Phase diagram of  $xy$ -trajectories:





## Related Literature (II)

- Economics and Quantum Game Theory

2002, E. W. Piotrowski and J. Sladkowski, *Quantum Market Games*, Physica A (312) 208

2002, Kay-Yut Chen, T. Hogg and R. Beaulsoleil *A Quantum Treatment of Public Goods Economics*, Quantum Information Processing 1(6)

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2007, T. Hogg, P. Harsha and Kay-Yut Chen *Quantum Auctions*, Int. J. of Quantum Information 5:751-780

2007, M. Hanauske, S. Bernius and B. Dugall, *Quantum Game Theory and Open Access Publishing*, Physica A, Vol.382 (2007), p.650-664 (physics/0612234)

## Related Literature (III)

- Quantum Computer and Quantum Game Theory

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*Experimental realization of quantum games on a quantum computer*, PRL 88 (137902)

2007, R. Prevedel, A. Stefanov, P. Walther and A. Zeilinger  
*Experimental realization of a quantum game on a one-way quantum computer*, New Journal of Physics 9 (205)

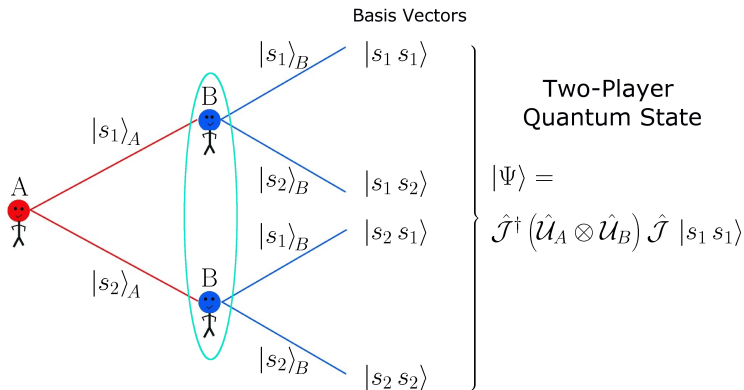
2008, P. Benicio, Melo de Sousa, R. V. Ramos *Multiplayer Quantum Games and its Application as Access Controller in Architecture of Quantum Computers*, arXiv:0802.3684v2

- Extensions of Quantum Game Theory

2001, S. C. Benjamin and P. M. Hayden, *Multi-Player Quantum Games*, PRA 64 (030301) ...



# A 2-Player-2-Strategy-Quantumgame



**Figure:**  $|\Psi\rangle$ : Two-Player State,  $\hat{\mathcal{J}}(\gamma)$ : Entangling Operator,  $\gamma$ : Strength of Entanglement,  $\hat{U}_A, \hat{U}_B$ : Strategy Decision Operator of Player A and B

# The Two-Player Quantum Wavefunction $|\Psi\rangle$

## The Two-Player Quantum State $|\Psi\rangle$

$$|\Psi\rangle = \hat{J}^\dagger (\hat{U}_A \otimes \hat{U}_B) \hat{J} |s_1 s_1\rangle$$

$\hat{U}_A$  : Decision Operator of Player A

$\hat{U}_B$  : Decision Operator of Player B

$\hat{J}$  : Entangling Operator

$\hat{J}^\dagger$  : Disentangling Operator

$\hat{J} |s_1 s_1\rangle$  : Two-Player Initial State ( $|\Psi_0\rangle$ )

## In words ...

The setup of the quantum game begins with the choice of the initial state  $|\Psi_0\rangle$ . After the two players have chosen their individual quantum strategies ( $\hat{U}_A := \hat{U}(\theta_A, \varphi_A)$  and  $\hat{U}_B := \hat{U}(\theta_B, \varphi_B)$ ) the disentangling operator  $\hat{J}^\dagger$  is acting to prepare the measurement.

# The quantum decision state $|\psi\rangle_\mu$ of player $\mu = A, B$

To illustrate the operator formalism of quantum game theory and the concept of quantum strategies, we want to focus at first on the real and imaginary values of the two spinor components  $\psi_1^A$  and  $\psi_2^A$  of the of the state  $|\psi\rangle_A$  of player A:

$$|\psi\rangle_A = \psi_1^A |s_1^A\rangle + \psi_2^A |s_2^A\rangle = \begin{pmatrix} \psi_1^A \\ -\psi_2^A \end{pmatrix} \in \mathcal{H}_A$$

$$|s_1^A\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \quad |s_2^A\rangle = \begin{pmatrix} 0 \\ -1 \end{pmatrix}$$

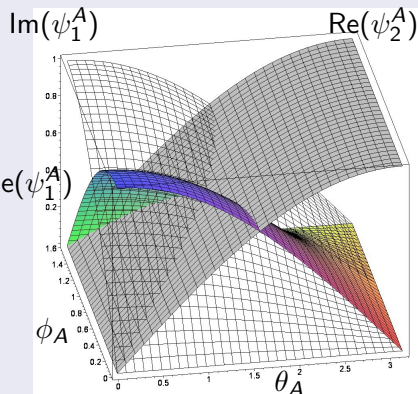
$$|\psi\rangle_A = \hat{U}(\theta_A, \varphi_A) |s_1^A\rangle = \begin{pmatrix} e^{i\varphi_A} \cos(\frac{\theta_A}{2}) \\ -\sin(\frac{\theta_A}{2}) \end{pmatrix}$$

$$\hat{U}(\theta, \varphi) := \begin{pmatrix} e^{i\varphi} \cos(\frac{\theta}{2}) & \sin(\frac{\theta}{2}) \\ -\sin(\frac{\theta}{2}) & e^{-i\varphi} \cos(\frac{\theta}{2}) \end{pmatrix} \quad \forall \theta \in [0, \pi] \wedge \varphi \in [0, \frac{\pi}{2}]$$



# The quantum decision state $|\psi\rangle_\mu$ of player $\mu = A, B$

Real and imaginary parts of  $|\psi\rangle_A$



Quantum state of player A:

$$|\psi\rangle_A = \psi_1^A |s_1^A\rangle + \psi_2^A |s_2^A\rangle = \begin{pmatrix} \psi_1^A \\ -\psi_2^A \end{pmatrix} \in \mathcal{H}_A$$

$$\text{with: } |s_1^A\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, |s_2^A\rangle = \begin{pmatrix} 0 \\ -1 \end{pmatrix}$$

$s_1$ -quantum strategies and the decision operator  $\hat{U}(\theta, \varphi)$ :

$$|\psi\rangle_A = \hat{U}(\theta_A, \varphi_A) |s_1^A\rangle = \begin{pmatrix} e^{i\varphi_A} \cos(\frac{\theta_A}{2}) \\ -\sin(\frac{\theta_A}{2}) \end{pmatrix}$$

$$\hat{U}(\theta, \varphi) := \begin{pmatrix} e^{i\varphi} \cos(\frac{\theta}{2}) & \sin(\frac{\theta}{2}) \\ -\sin(\frac{\theta}{2}) & e^{-i\varphi} \cos(\frac{\theta}{2}) \end{pmatrix}$$

$$\forall \theta \in [0, \pi] \wedge \varphi \in [0, \frac{\pi}{2}]$$

# Interpretation

Die quantentheoretische Beschreibung des Entscheidungszustandes des Spielers A kurz vor der definitiven Auswahl und Bekundung der reinen Strategie besitzt demnach im Allgemeinen neben den reelwertigen auch imaginäre Anteile. Bei  $s_1$ -Quantenstrategien kann sich der Spieler nur im imaginären Raum der ersten Strategie gedanklich bewegen. Eine grundlegende Eigenschaft der gesamten Quantentheorie ist die prinzipielle Unbeobachtbarkeit des Quantenzustandes. Diese Eigenschaft spiegelt sich in der Quanten Spieltheorie in der Unbeobachtbarkeit des Gedankenprozesses wider. Die einzelnen Gedankenwege, die während des Entscheidungsprozesses im Gehirn des Spielers (bewusst oder unterbewusst) ablaufenden, können nicht direkt gemessen werden.

$s_1$ -Quantenstrategien können als der gedankliche Weg während des Entscheidungsprozesses interpretiert werden, welcher vom gedanklichen Ursprung her von der klassischen Strategie  $s_1$  startet und hypothetisch, gebunden an die Wünsche und Ängste des Spielers, den Gedankenweg weiterbildet.

# The 2-player state $|\Psi\rangle$ and the entangling operator $\hat{\mathcal{J}}(\gamma)$

$$|\Psi\rangle = \hat{\mathcal{J}}^\dagger (\hat{u}_A \otimes \hat{u}_B) \hat{\mathcal{J}} |s_1^A s_1^B\rangle$$

$$\hat{\mathcal{J}} := e^{i\frac{\gamma}{2}(\hat{s}_1 \otimes \hat{s}_1)} = \begin{pmatrix} \cos\left(\frac{\gamma}{2}\right) & 0 & 0 & i \sin\left(\frac{\gamma}{2}\right) \\ 0 & \cos\left(\frac{\gamma}{2}\right) & -i \sin\left(\frac{\gamma}{2}\right) & 0 \\ 0 & -i \sin\left(\frac{\gamma}{2}\right) & \cos\left(\frac{\gamma}{2}\right) & 0 \\ i \sin\left(\frac{\gamma}{2}\right) & 0 & 0 & \cos\left(\frac{\gamma}{2}\right) \end{pmatrix}$$

$$\gamma \in [0, \frac{\pi}{2}], \quad |s_1^A s_1^B\rangle := |s_1^A\rangle \otimes |s_1^B\rangle = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

# Interpretation

The most important, but also most difficult mathematical concept in QGT is the two player quantum state  $|\Psi\rangle$ . It is formally constructed with the use of the decision operators  $\hat{U}_A$  and  $\hat{U}_B$  of player  $A$  and  $B$  and the entangling and disentangling operator  $\hat{J}$  and  $\hat{J}^\dagger$ .  $|\Psi\rangle$  is a spinor in a complex valued, 4-dimensional, abstract mathematical space called the 2-player "Hilbertspace"  $\mathcal{H}$ . The space of all conceivable decision paths is extended from the purely rational, measurable space in the Hilbertspace of complex numbers. Through the concept of a potential entanglement of the imaginary quantum strategy parts, it is possible to include corporate decision paths, caused by cultural or moral standards. QGT is therefore a model which goes beyond *Homo Economicus* and the parameter  $\gamma$ , which is a measure for the strength of entanglement and fellow feeling, describes how strongly the players behave as *Homo Sociologicus* or *Homo Transzendentalis*.

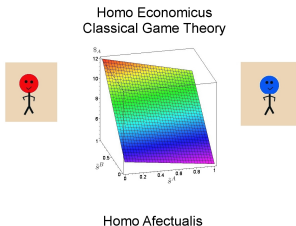
# The 2-player state $|\Psi\rangle$ and the entangling operator $\hat{\mathcal{T}}(\gamma)$

## Beyond Homo Economicus

Quantum Game Theory  
Entanglement    Quantum Strategies  
Homo Sociogenicus    Homo Transcendentalis

$$|\Psi\rangle$$

Extended models of classical evolutionary game theory (e.g. [10, 9])



The final 2-player quantum state:

$$|\Psi\rangle = \hat{\mathcal{T}}^\dagger \left( \hat{U}_A \otimes \hat{U}_B \right) \hat{\mathcal{T}} |s_1^A s_1^B\rangle$$

$\hat{\mathcal{T}}(\gamma)$ : Entangling operator

$\hat{\mathcal{T}}^\dagger(\gamma)$ : Disentangling operator

$\gamma \in [0, \pi]$ : Strength of entanglement

$\hat{U}_A$ : Decision Operator for player A

$\hat{U}_B$ : Decision Operator for player B

# The extended payoff $\$_{\mu}(\theta_A, \varphi_A, \theta_B, \varphi_B)$ of player $\mu = A, B$

The extended payoff  $\$_{\mu}(\tau_A, \tau_B)$  of player  $\mu = A, B$  is an amplification of the classical mixed strategy payoff function  $\tilde{\$}^{\mu}(\tilde{s}^A, \tilde{s}^B)$ :

$$\$A = \$_{11}^A P_{11} + \$_{12}^A P_{12} + \$_{21}^A P_{21} + \$_{22}^A P_{22}$$

$$\$B = \$_{11}^B P_{11} + \$_{12}^B P_{12} + \$_{21}^B P_{21} + \$_{22}^B P_{22}$$

$$\text{with: } P_{\sigma\sigma'} = |\langle \sigma\sigma' | \Psi \rangle|^2, \quad \sigma = \{s_1^A, s_2^A\} \text{ and } \sigma' = \{s_1^B, s_2^B\}$$

$P_{\sigma\sigma'}$  are the real valued probabilities of finding the two player state  $|\Psi\rangle$  in the pure strategy Eigenstate  $|\sigma\sigma'\rangle$ , e.g.

$$P_{12} := P_{s_1^A s_2^B} = \left| \langle s_1^A s_2^B | \Psi \rangle \right|^2$$

## The extended payoff $\$_{\mu}(\tau_A, \tau_B)$ of player $\mu = A, B$

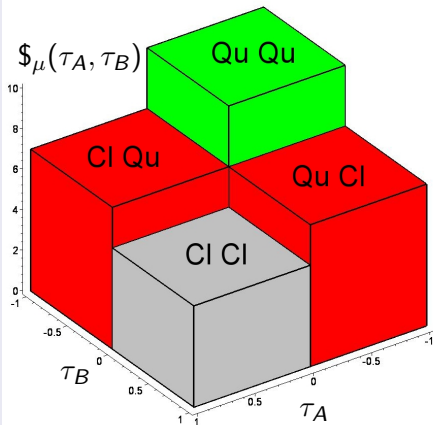
In contrast to the classical mixed payoff functions ( $\tilde{\$}^A(\tilde{s}^A, \tilde{s}^B)$  and  $\tilde{\$}^B(\tilde{s}^A, \tilde{s}^B)$ ), which depend only on the two parameters  $\tilde{s}^A$  and  $\tilde{s}^B$ , the quantum version of the mixed strategy payoff function depends in general on five parameters; namely the four decision angles of player A and player B ( $\theta_A, \varphi_A, \theta_B$  and  $\varphi_B$ ) and the entangling parameter  $\gamma$ . In order to visualize the payoff function as a surface in a three dimensional space it is necessary to reduce the set of parameters in the final state:

$|\Psi\rangle = |\Psi_f(\theta_A, \varphi_A, \theta_B, \varphi_B)\rangle \rightarrow |\Psi(\tau_A, \tau_B)\rangle$ . The two strategy angles  $\theta$  and  $\varphi$  depend only on a single parameter  $\tau \in [-1, 1]$ .

Positive  $\tau$ -values represent pure and mixed classical strategies, whereas negative  $\tau$ -values correspond to quantum strategies, where  $\theta = 0$  and  $\varphi > 0$ . The whole strategy space is separated into four regions, namely the absolute classical region (*ClCl*:  $\tau_A, \tau_B \geq 0$ ), the absolute quantum region (*QuQu*:  $\tau_A, \tau_B < 0$ ) and the two partially classical-quantum regions (*ClQu*:  $\tau_A \geq 0 \wedge \tau_B < 0$  and *QuCl*:  $\tau_A < 0 \wedge \tau_B \geq 0$ ).

# The extended payoff $\$_{\mu}(\tau_A, \tau_B)$ of player $\mu = A, B$

## Visualisationspace of $\$_{\mu}(\tau_A, \tau_B)$



The expected payoff within a quantum version of a general 2-player game:

$$\$A = \$_{11}^A P_{11} + \$_{12}^A P_{12} + \$_{21}^A P_{21} + \$_{22}^A P_{22}$$

$$\$B = \$_{11}^B P_{11} + \$_{12}^B P_{12} + \$_{21}^B P_{21} + \$_{22}^B P_{22}$$

$$\text{with: } P_{\sigma\sigma'} = |\langle \sigma\sigma' | \Psi \rangle|^2, \quad \sigma, \sigma' = \{s_1, s_2\}$$

Reduction of quantum strategies:

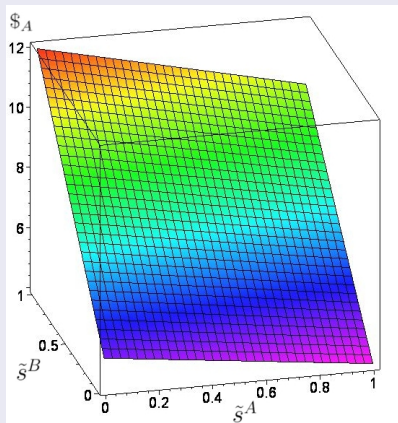
$$|\Psi\rangle = |\Psi(\theta_A, \varphi_A, \theta_B, \varphi_B)\rangle \rightarrow |\Psi(\tau_A, \tau_B)\rangle$$

$$\underbrace{\{(\tau, 0) \mid \tau \in [0, 1]\}}_{\text{classical region } Cl} \wedge \underbrace{\{(0, \tau) \mid \tau \in [-1, 0]\}}_{\text{quantum region } Qu}$$



# Quantum extension of dominant class games

## Classical payoff for player A



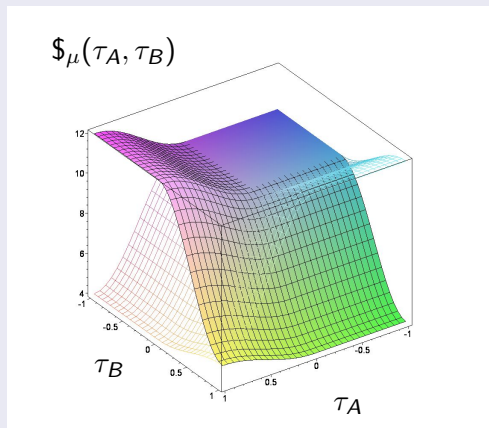
A \ B	$s_1^B$	$s_2^B$
$s_1^A$	(10,10)	(4,12)
$s_2^A$	(12,4)	(5,5)

**Table:** Payoffmatrix of a dominant, prisoners dilemma like game.

This dominant, prisoners dilemma like game has only one pure, symmetric Nash equilibrium  $(s_2^A, s_2^B)$  which is the only ESS of the evolutionary game.

# Quantum extension of dominant class games

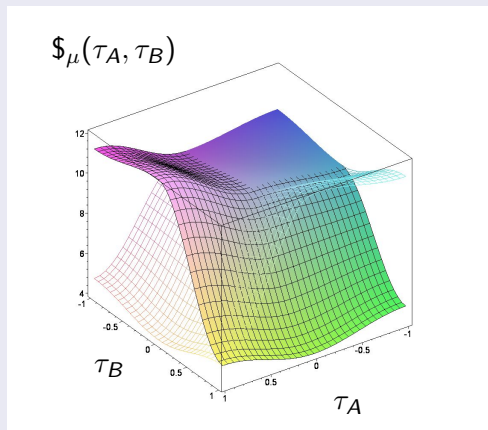
Payoff of player A (colored) and player B (wired) for  $\gamma = 0$  (no entanglement)



The diagram clearly exhibits that the non-entangled quantum game simply describes the classical version of the prisoner's dilemma game. For the case, that both players decide to play a quantum strategy ( $\tau_A < 0 \wedge \tau_B < 0$ ) their payoff is equal to the case where both players choose the classical pure strategy  $s_1$  ( $\$_A(\tau_A = 0, \tau_B = 0) = 10$ ). The classical Nash equilibrium ( $(s_2^A, s_2^B)$ , the dominant strategy) corresponds to the following  $\tau$ -values:  $(s_2^A, s_2^B) \hat{=} (\tau_A = 1, \tau_B = 1)$ .

# Quantum extension of dominant class games

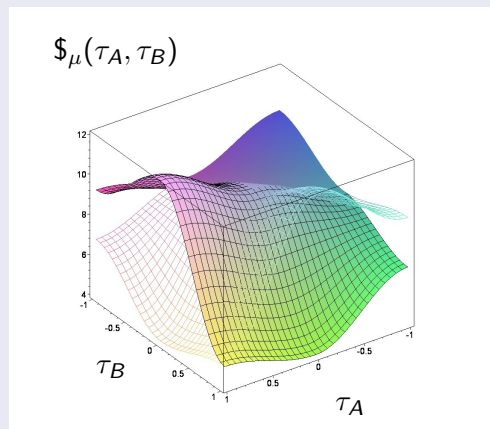
Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{10} \approx 0.31$



For the absolute classical region *CICl* the shape of the surfaces does not change, whereas for the partially classical-quantum (*ClQu* and *QuCl*) and absolute quantum region regions *QuQu* the payoff structure changes due to a possible interference of quantum strategies within Hilbertspace. The structure of Nash-equilibria does not change for the left picture, whereas for the following pictures the previously present dominant strategy of the prisoner's dilemma game has disappeared and a new, advisable quantum Nash-equilibrium will appear at  $(\hat{Q}, \hat{Q} \hat{=} (\tau_A = -1, \tau_B = -1))$ . During the transition from this figure to the next picture two separate phenomena occur. At first, for an entanglement value  $\gamma_1 \approx 0.37$ , the best response for player A to the strategy  $s_2^B \hat{=} \tau_B = 1$  is no longer the strategy  $s_2^A \hat{=} \tau_A = 1$ , as  $\$_A(\tau_A = -1, \tau_B = 1) \approx 5.05$  is now higher than  $\$_A(\tau_A = 1, \tau_B = 1) = 5$ . Secondly, for an entanglement value  $\gamma_2 \approx 0.53$ , the best response for player A to the strategy  $\hat{Q}_B \hat{=} \tau_B = -1$  is no longer the strategy  $s_2^A \hat{=} \tau_A = 1$ , as  $\$_A(\tau_A = 1, \tau_B = -1) \approx 9.96$  is for  $\gamma_2 = 0.53$  lower than  $\$_A(\tau_A = -1, \tau_B = -1) = 10$ .

# Quantum extension of dominant class games

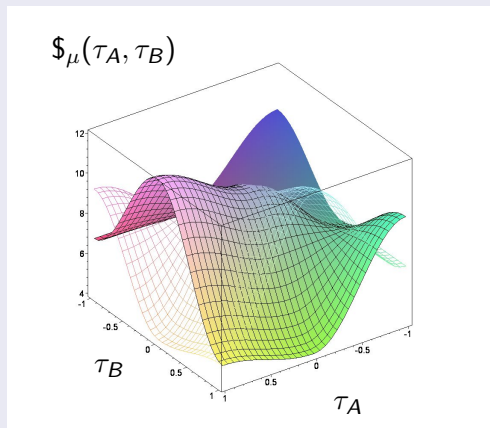
Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{8} \approx 0.52$



For the absolute classical region *CICl* the shape of the surfaces does not change, whereas for the partially classical-quantum (*ClQu* and *QuCl*) and absolute quantum region regions *QuQu* the payoff structure changes due to a possible interference of quantum strategies within Hilbertspace. The structure of Nash-equilibria did not change for the last figure, whereas for this and the following pictures the previously present dominant strategy of the prisoner's dilemma game has disappeared and a new, advisable quantum Nash-equilibrium has appeared ( $\hat{Q}, \hat{Q} \hat{=} (\tau_A = -1, \tau_B = -1)$ ). During the transition from the last picture to this figure two separate phenomena occurred. At first, for an entanglement value  $\gamma_1 \approx 0.37$ , the best response for player A to the strategy  $s_2^B \hat{=} \tau_B = 1$  is no longer the strategy  $s_2^A \hat{=} \tau_A = 1$ , as  $\$_A(\tau_A = -1, \tau_B = 1) \approx 5.05$  is now higher than  $\$_A(\tau_A = 1, \tau_B = 1) = 5$ . Secondly, for an entanglement value  $\gamma_2 \approx 0.53$ , the best response for player A to the strategy  $\hat{Q}_B \hat{=} \tau_B = -1$  is no longer the strategy  $s_2^A \hat{=} \tau_A = 1$ , as  $\$_A(\tau_A = 1, \tau_B = -1) \approx 9.96$  is for  $\gamma_2 = 0.53$  lower than  $\$_A(\tau_A = -1, \tau_B = -1) = 10$ .

# Quantum extension of dominant class games

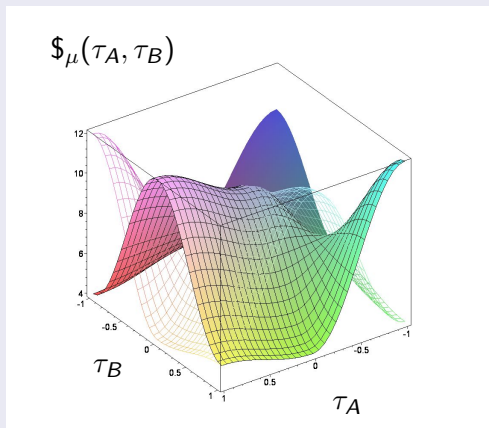
Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{6} \approx 0.94$



The results show, that a quantum extension of a classical prisoner's dilemma game is able to change the structure of Nash-equilibria, and even previously present dominant strategies could become nonexistent, if the value of entanglement increases further than a defined  $\gamma$ -threshold. Players with a higher strategic entanglement value  $\gamma$  escape the dilemma as they see the advantage of the quantum strategy combination  $(\hat{Q}_A, \hat{Q}_B)$ , which is measured as if both are playing the classical strategy  $s_2$ .

# Quantum extension of dominant class games

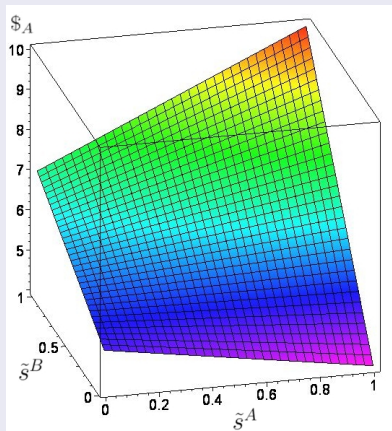
Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{2} \approx 1.57$



The results show, that a quantum extension of a classical prisoner's dilemma game is able to change the structure of Nash-equilibria, and even previously present dominant strategies could become nonexistent, if the value of entanglement increases further than a defined  $\gamma$ -threshold. Players with a higher strategic entanglement value  $\gamma$  escape the dilemma as they see the advantage of the quantum strategy combination  $(\hat{Q}_A, \hat{Q}_B)$ , which is measured as if both are playing the classical strategy  $s_2$ .

# Quantum extension of coordination class games

## Classical payoff for player A



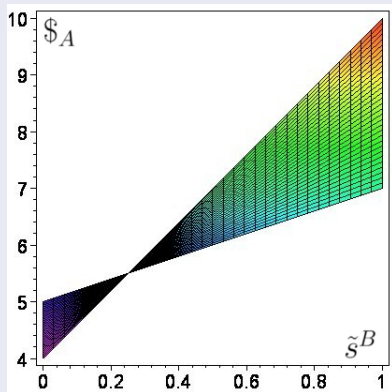
A \ B	$s_1^B$	$s_2^B$
$s_1^A$	(10,10)	(4,7)
$s_2^A$	(7,4)	(5,5)

**Table:** Payoffmatrix of a coordination game.

This coordination game has two pure, symmetric Nash equilibria and one interior NE at  $s^* = \frac{1}{4}$ . The evolutionary game game has two ESSs.

# Quantum extension of coordination class games

## Classical payoff for player A (projected)



A \ B	$s_1^B$	$s_2^B$
$s_1^A$	(10,10)	(4,7)
$s_2^A$	(7,4)	(5,5)

**Table:** Payoffmatrix of a coordination game.

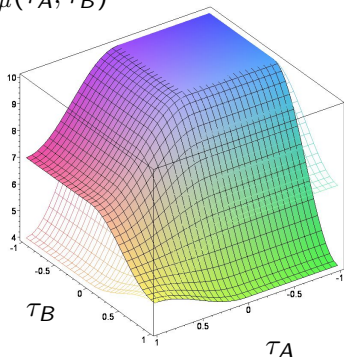
This coordination game has two pure, symmetric Nash equilibria and one interior NE at  $s^* = \frac{1}{4}$ . The evolutionary game game has two ESSs.



# Quantum extension of coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = 0$  (no entanglement)

$$\$_{\mu}(\tau_A, \tau_B)$$

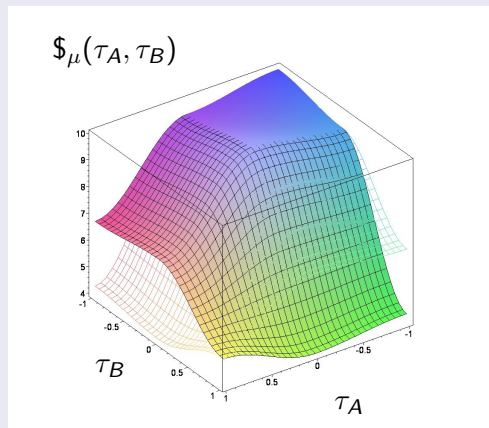


Again, the diagram clearly indicates that the non-entangled quantum game is identical to the classical version of the underlying coordination game. For the case, that both players decide to play a quantum strategy ( $\tau_A < 0 \wedge \tau_B < 0$ ) their payoff is equal to the case where both players choose the classical pure strategy  $s_1$  ( $\$A(\tau_A = 0, \tau_B = 0) = 10$ ), with the overall highest possible payoff. The classical pure Nash equilibria correspond to the following  $\tau$ -values:  $(s_1^A, s_1^B) \hat{=} (\tau_A = 0, \tau_B = 0)$  and  $(s_2^A, s_2^B) \hat{=} (\tau_A = 1, \tau_B = 1)$ , whereas the classical mixed strategy equilibrium is at:

$$\tau^* = \frac{2}{\pi} \arccos(\sqrt{\frac{1}{4}}) = \frac{2}{3}.$$

# Quantum extension of coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{10} \approx 0.31$



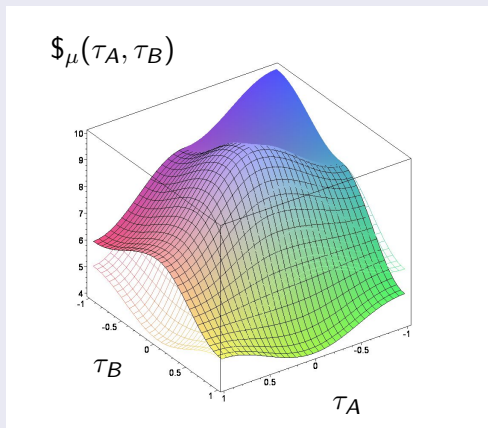
Even for tiny values of  $\gamma$  a new quantum Nash-equilibrium appears ( $\tau_A = -1, \tau_B = -1$ ).

At moderate values of  $\gamma$  the low payoff evolutionary stable strategy ( $\tau_A = 1, \tau_B = 1$ ) disappears.

The specific  $\gamma$ -value at which this disappearance happens, depends on the whole set of payoff parameters and not only on  $a$  and  $b$ .

# Quantum extension of coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{8} \approx 0.52$



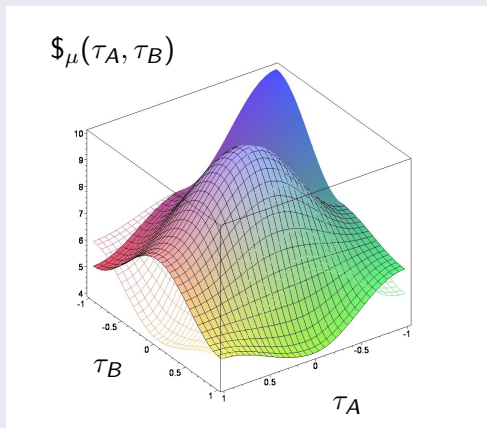
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# Quantum extension of coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{6} \approx 0.94$



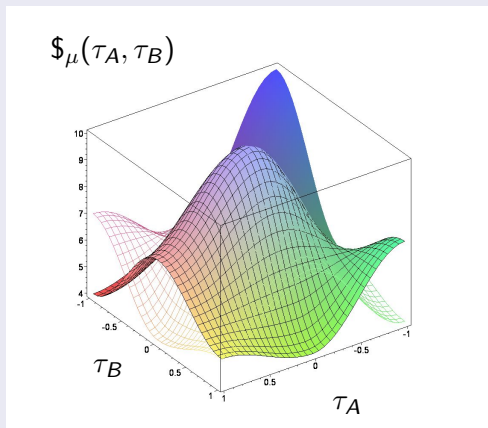
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# Quantum extension of coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{2} \approx 1.57$



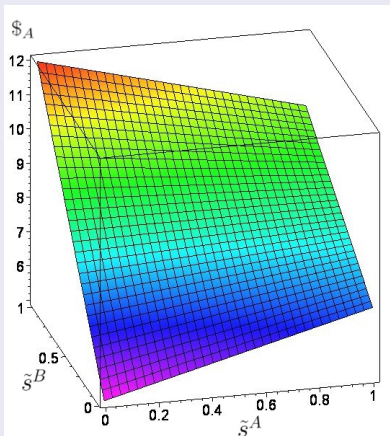
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# Quantum extension of anti-coordination class games

## Classical payoff for player A



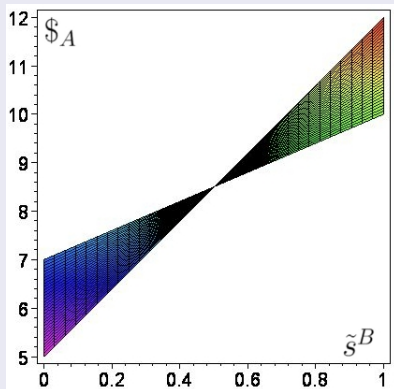
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Table: Payoffmatrix of a coordination game.

This anti-coordination game has two pure, unsymmetric Nash equilibria and one interior NE at  $s^* = \frac{1}{2}$ . The evolutionary game game has one mixed strategy ESS.

# Quantum extension of anti-coordination class games

## Classical payoff for player A (projected)



A \ B	$s_1^B$	$s_2^B$
$s_1^A$	(10,10)	(7,12)
$s_2^A$	(12,7)	(5,5)

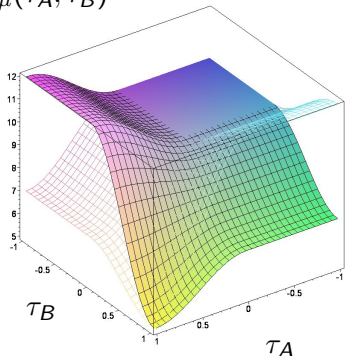
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# Quantum extension of anti-coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = 0$

$$\$_\mu(\tau_A, \tau_B)$$



Beside the mixed strategy evolutionary stable strategy, a new quantum ESS appears at a specific  $\gamma$ -value.

For details see:

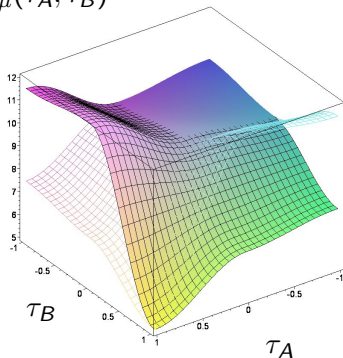
- M. Hanauske, *Advances in Evolutionary Game Theory*, 2009, Lecture at the 'Université Lumière Lyon 2' in Lyon, France (MINERVE Exchange Program); Slides and additional material
- M. Hanauske, J. Kunz, S. Bernius, and W. König, 'Doves and hawks in economics revisited: An evolutionary quantum game theory-based analysis of financial crises.', 2009, to appear in *Physica A*, arXiv:0904.2113, RePEc:pra:mprapa:14680 and SSRN<sub>id</sub>:1597735.



# Quantum extension of anti-coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{10} \approx 0.31$

$$S_{\mu}(\tau_A, \tau_B)$$



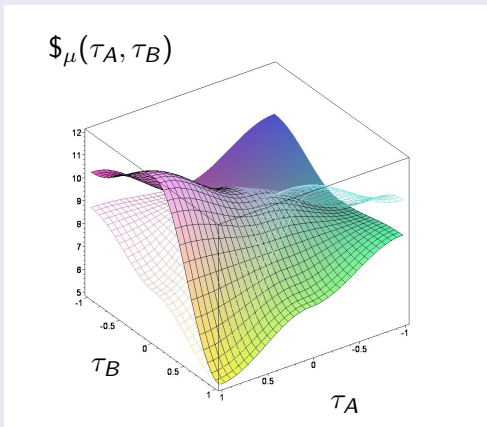
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# Quantum extension of anti-coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{8} \approx 0.52$



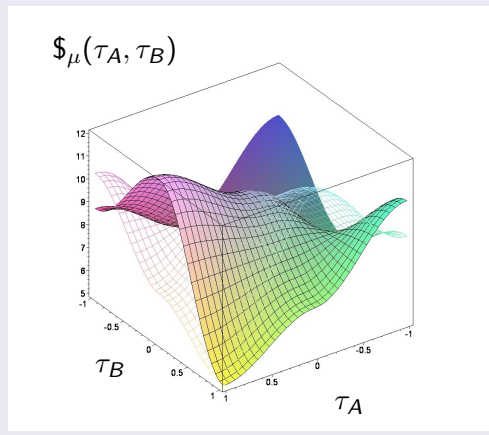
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# Quantum extension of anti-coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{6} \approx 0.94$



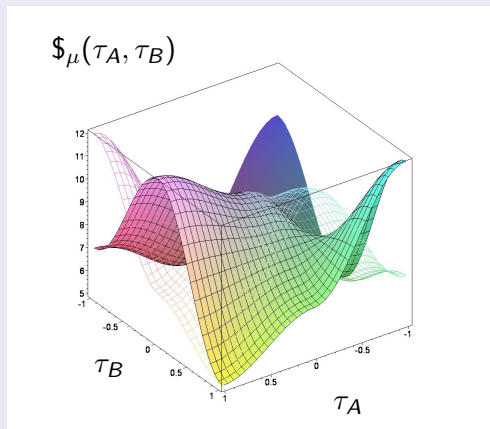
Beside the mixed strategy evolutionary stable strategy, a new quantum ESS appears at a specific  $\gamma$ -value.

For details see:

- M. Hanauske, *Advances in Evolutionary Game Theory*, 2009, Lecture at the 'Université Lumière Lyon 2' in Lyon, France (MINERVE Exchange Program); Slides and additional material
- M. Hanauske, J. Kunz, S. Bernius, and W. König, 'Doves and hawks in economics revisited: An evolutionary quantum game theory-based analysis of financial crises', 2009, to appear in *Physica A*, arXiv:0904.2113, RePEc:pra:mpnpra:14680 and SSRN<sub>id</sub>:1597735.

# Quantum extension of anti-coordination class games

Payoff of player A (colored) and player B (wired) for  $\gamma = \frac{\pi}{2} \approx 1.57$



Beside the mixed strategy evolutionary stable strategy, a new quantum ESS appears at a specific  $\gamma$ -value.

For details see:

- M. Hanauske, *Advances in Evolutionary Game Theory*, 2009, Lecture at the 'Université Lumière Lyon 2' in Lyon, France (MINERVE Exchange Program); Slides and additional material
- M. Hanauske, J. Kunz, S. Bernius, and W. König. ,Doves and hawks in economics revisited: An evolutionary quantum game theory-based analysis of financial crises. , 2009, to appear in *Physica A*, arXiv:0904.2113, RePEc:pra:mprapa:14680 and SSRN<sub>id</sub>:1597735 .

# Quantum Game Theory and Scientific Communication

M. Hanauske, S. Bernius, and B. Dugall, *Quantum Game Theory and Open Access Publishing, Physica A*, 382(2):650–664, 2007, arXiv:physics/0612234

Matthias Hanauske, Wolfgang König, and Berndt Dugall, *Evolutionary Quantum Game Theory and Scientific Communication*, 2010, Accepted article of the "Second Brazilian Workshop of the Game Theory Society", Internet-Link

Matthias Hanauske, *Evolutionary Game Theory and Complex Networks of Scientific Information*, Article is going to be published as a chapter in the book "*Models of science dynamics - Encounters between complexity theory and information science*", Springer book in the Complexity series, Editors: Andrea Scharnhorst, Katy Börner and Peter van den Besselaar, 2011

# Quantum Game Theory and Scientific Communication

In recent years the market of scientific publishing faces several forces that may cause a major change of traditional market mechanisms. In the following we understand open access publishing as the electronic publication of scientific information on a platform that provides access to this information for all potential users, without financial or other barriers. In contrast most other scientific disciplines do not make use of open access publishing, even though they support this model if asked for [2]. Instead, they submit research papers to traditional journals that do not provide free access to their articles. Considering that the majority of scientists regard open access publishing as superior to the traditional system, the question arises, why it is only adopted by few disciplines.



# Quantum Game Theory and Scientific Communication

The payoff structure of this game is modeled by the following payoff matrix:

A \ B	o	∅
o	$(r + \delta, r + \delta)$	$(r - \alpha, r + \beta)$
∅	$(r + \beta, r - \alpha)$	$(r, r)$

**Table:** Researchers open access payoff matrix.

The actual reputation of the two scientists is represented by a single parameter  $r$ . If only one of the two players chooses the open access strategy ( $(\emptyset, o)$  or  $(o, \emptyset)$ ) the parameters  $\alpha$  and  $\beta$  ( $\alpha, \beta \geq 0$ ) describe the decrease and the increase of the scientists' reputation, depending on the selected strategy. The parameter  $\delta$  represents the potential benefit in the case that both players choose the open access strategy  $(o, o)$ .







# Summary

## Summary of the talk

Quantum game theory is a mathematical and conceptual amplification of classical game theory. The space of all conceivable decision paths is extended from the purely rational, measurable space in the Hilbertspace of complex numbers. Through the concept of a potential entanglement of the imaginary quantum strategy parts, it is possible to include corporate decision path, caused by cultural or moral standards. If this strategy entanglement is large enough, then, additional Nash-equilibria can occur, previously present dominant strategies could become nonexistent and new evolutionary stable strategies can appear.

Within this talk the framework of Quantum Game Theory was described in detail. The formal mathematical model, the different concepts of equilibria and the various classes of quantum games have been defined, explained and visualized to understand the main ideas of Quantum Game Theory. Additionally, two applications were discussed at the end of the talk.









J.W. Weibull.

*Evolutionary Game Theory.*

The MIT Press, 1995.